

CSS 342 - Mathematical Principles of Computing

- Data structures, maintainable, good quality code
- dude uses a scripter to grade → strict on naming, style conventions
- 5 projects → start early! C++ has pitfalls
- char strings end in "0" ⇒ Test code outside of Visual Studio
- zero out memory contents or face unpleasant consequences
- globals vs locals } use getters
- Interpreter vs. Compiler
- Structs require semi-colons
- ints have different bit sizes in C++
- Use message board for 332 or even 342 Questions
- email Nash if needed

Clearing transaction data → memory idempotent

Shields uses G++ compiler

constants in header files

CSS 342 C: Mathematical Principles of Computing p.4/4				
Canvas/Catalyst EPost (Message Board): You can access the message board link on the class website to exchange messages with your classmates. Please use this board for only discussions. No junk email. Note that the professor will not keep track of all messages. The professor monitors this board once a day Mon. & Wed.				
Topics covered and tentative schedule: This is an approximate ordering of topics. Material will take about the allotted time; we do not cover all sections in all chapters. See Website for updates. Readings prefixed with C++ or Math correspond to the appropriate textbook.				
week	Date	Topic (topics and work are subject to change by instructor as needed)	Read this Material before class 6th ed. chapters for each book are in ()	Lab(Prog-ram) Due
0	Sep 25	Introduction, etc. C++ language structure, C++ structs, arrays	Website C++ Notes, C++ appendix A, G,H,I,K,J as needed	
1	Sep 30	Software Engineering Principles, pair programming; Searching, simple note: p.71 (L-f)/2sorting (selection, bubble, insertion)	C++ 1, 9.2 (6ed. appendices BCDF, 2.1, 2.4.2, 10.2.5, 11.1, interlude 3 for asset, appx I for doxygen, G for file I/O, L for library includes), Website Notes	
	Oct 2	Data Abstraction (Objects and Classes), C++ classes	Website C++ Notes C++ 3 (6ed. 1, Interlude 1))	
2	7	More C++ classes	"	lab 1 due
	9	Pointers, Linked Lists	C++ 4.1-4.4 (6ed. Interlude2; ch.4,8,9,14.1.2)	
3	14	more Pointers, Linked	Website	lab 2 due
	16	Templates, STL intro	C++ 4.5, apx E (6ed. 1.8.3, Interlude 1 & 7)	
4	21	Algorithm analysis	C++ 9.1 (6ed. 10, 11.3) ?Math 2.4, 3.1-3.3	
	23	Algorithm analysis continued		
5	28	Recursion	C++ 2 (6ed. 2, 3.3, 4.3, 5, 6.5, 9.2.3) ?Math 5.3-5.4 (4.3-4.4 ed. 6)	lab 3 due
	30	Midterm exam		
6	Nov 4	Sorting (using recursion)	C++ 9.2 (some of) (6ed. 11)	
	6	Binary Search Tree intro	C++ 10.3 (6ed. 15, 16), Website links/notes ?Math 11.1-11.3 (10.1-10.3 ed. 6)	
7	11	No school - holiday		
	12	Last day to drop a course (& a late fee is assessed)		
	13	Object-oriented design, Stacks	C++ 1, 7 (6ed. 1, 6, 7)	lab 4 due
8	18	Queues	C++ 7 (6ed. 13,14)	
	20	Recurrence and Induction	C++ 2.1, 5.3 (6ed. 2.2,2.3, 5.4, appendix E), ?Math 5.1-5.2 (4.1-4.2 ed. 6)	lab 5 design due
9	25	Recurrence and Induction continued	C++ appendix D ?Math 8.1, 8.3 (7.1, 7.3 ed. 6)	
	27	Representation of Integers	?Math 4.2 (3.6 ed. 6) Website links/notes	
10	Dec 2	Mathematical Foundation	?Math 1.1, 1.3-1.5 (1.1-1.4 ed. 6)	
	4	Intro to Design Patterns, last day stuff	Website Course Notes	lab 5 due
11	9	Final exam, Monday, in class		

Appendix G - File I/O → Read

doxygen - page 777

H - Some standard headers

K - From Java to C++

Object-Oriented Solution

use classes → cohesive modules and loosely coupled modules
e.g. read module, sort module, print module

Operation contract → docs use specify pre- and post-conditions
review slides

Abstract Data Type - isolated inwards



Note: Notice that the C++ computation of the midpoint `mid` is

```
int mid = first + (last - first) / 2;
```

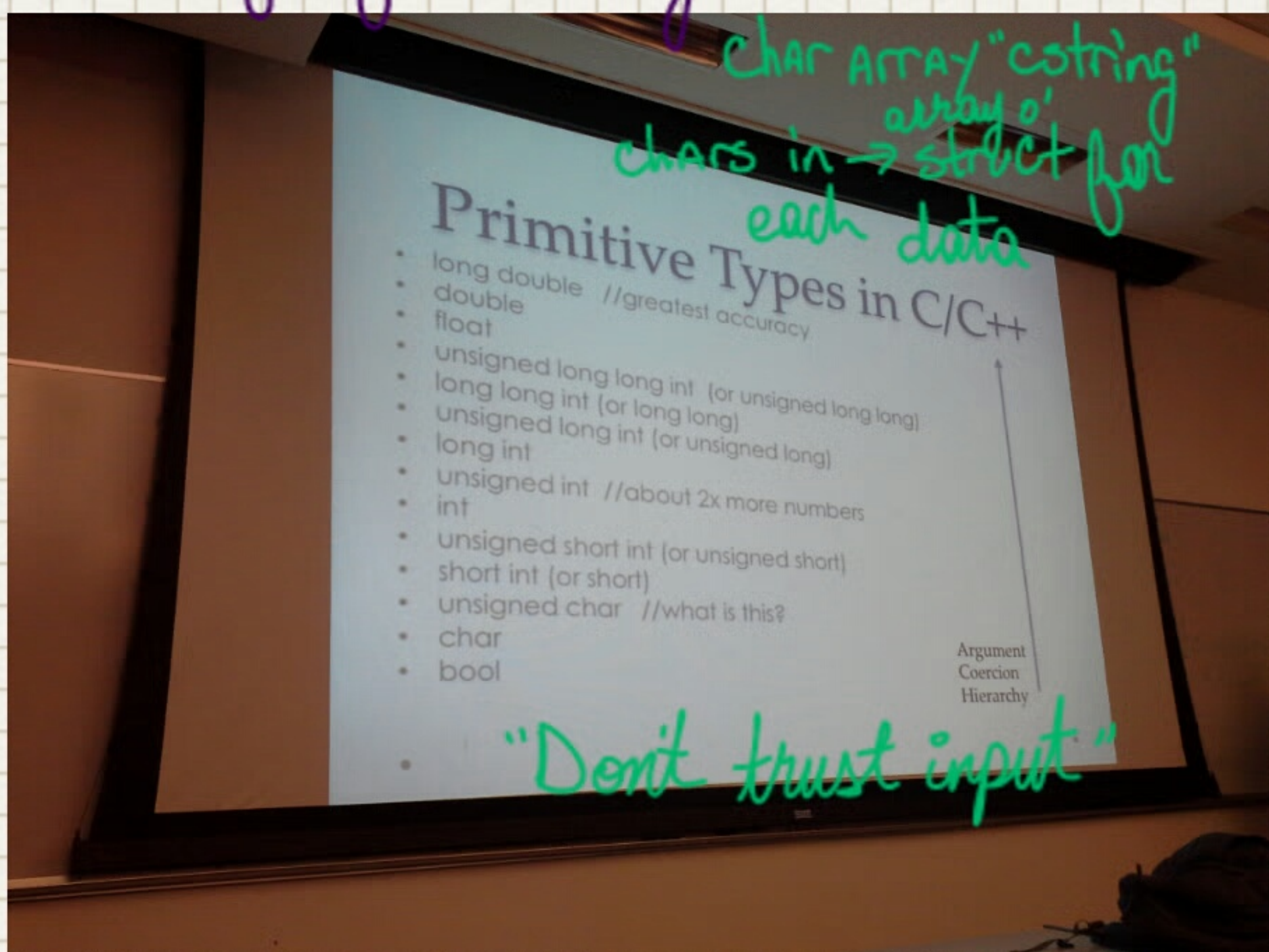
instead of

```
int mid = (first + last) / 2;
```

↑ I don't see how these are equal

as the pseudocode would suggest. If you were to search an array of at least 2^{30} , or about 1 billion, elements, the sum of `first` and `last` could exceed the largest possible `int` value of $2^{30} - 1$. Thus, the computation `first + last` would overflow to a negative integer and result in a negative value for `mid`. If this negative value of `mid` was used as an array index, it would be out of bounds and cause incorrect results. The computation `first + (last - first) / 2` is algebraically equivalent to `(first + last) / 2` and avoids this error.

Light discussion of sorting → selection, binary, bubble → see 143 notes
Utility functions (Decompose project #1)
structs for first assignment. headers declare mains dependency divide & conquer



When passing in char array
use an array of structs
use string copy

Look up Const

Pointers to Primitives

- `int c=0;` //addr 0x300, val is 0
- `int* cPtr = &c;`
- `int b;` //addr is 0x500, val is ?
- `int* bPtr = &b;`

Line of Code	c	cPtr	b	b_Ptr
(after above)	0	0x300	garbage	0x500
<code>c=5;</code>	?			
<code>*cPtr=3;</code>	?			
<code>cPtr = bPtr;</code>		?		
<code>*cPtr = 20;</code>	?		?	
<code>cPtr = null;</code>	?	?	?	?

Pointers to Arrays

- Given an int a, such as...
- `int a = 3;`
- `const int SIZE = 30;`
- `int* aPtr = &a;`
- `int a[] = {1,2,3};` //or
- `int* a = new int[3];` `a[0]=1, a[1]=2`
- `int*& aRef = a;` //what is this?
- `const int* aPtr = &SIZE;` //would &a work here?
- `int* const bPtr = &a;` //can I change this pointer?

9 7 5 3 1 0
x 9 7 5 3 1 0

9 7 5 3 1 0
7 9 3 5 0 1
0 1 3 5 7 9

11, 10, 9, 8, 7, 6, 5, 4, 3, 2, 1, 0 $m=12$

[11, 10, 9, 8, 7, 6] [5, 4, 3, 2, 1, 0] ①

[11, 10, 9] [8, 7, 6] [5, 4, 3] [2, 1, 0] ②

[11, 10] [9] [8, 7] [6] [5, 4] [3] [2, 1] [0] ③

[11] [10] [9] [8] [7] [6] [5] [4] [3] [2] [1] [0] ④

[11, 10, 9, 8, 7, 6, 5, 4, 3, 2, 1, 0] = A_0

[0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0] = B_0 $0 = \text{make } B$

[11, 10, 9, 8, 7, 6, 0, 0, 0, 0, 0, 0] = B_1

[0, 0, 0, 0, 0, 0, 5, 4, 3, 2, 1, 0] = A_1 $1 = A \rightarrow B$

[0, 0, 0, 8, 7, 6, 0, 0, 0, 0, 0, 0] = B_2

[11, 10, 9, 0, 0, 0, 5, 4, 3, 2, 1, 0] = A_2 $2 = B \rightarrow A$

[11, 10, 0, 8, 7, 6, 0, 0, 0, 0, 0, 0] = B_3 $4 = B \rightarrow A$

[0, 0, 9, 0, 0, 0, 5, 4, 3, 2, 1, 0] = A_3 $3 = A \rightarrow B$

[11, 0, 9, 0, 0, 0, 5, 4, 3, 2, 1, 0] = A_4

[0 10 0

Object-Oriented Concepts

- Code using a solution design
- Specify a system of interacting objects
- Object-oriented *analysis* specifies
 - What to do
 - Not how to do it

Review
Bag ADT

Top Down Design!

- Object-oriented *design* specifies
 - Models of how it might be done

Abstract out sub procedures

subprocedures
- Get DATA
- Insert

Data Structures and Problem Solving with C++ Walls and Mirrors, Garrano and Henry, © 2013

Steps to develop a solution:

1. Object Oriented Analysis → "Discovery" of problem. Not focusing on soln

- What does the end user expect?
- Requirements → what a soln must be and do.
- Express the problem and requirements of the soln in terms of objects
- Describe these objects and their interactions

^{Default} Constructor broke equality → didn't initialize

Interlude 1: classes

- If a method does not change an object, mark it as const
 - source in a copy method.

p.40 → Inheritance

plaintext examples

Linked lists in C++ always lives on heap

Dynamic data structures live on heap

Arrays can hold primitives and references

Make our own better versions
→ Bounds checking, deep copies
Automatic size mgmt, dynamic

Multi-Dimensional Arrays

- `int arr[3][3] = { {1,2,3}, {4,5,6}, {7,8,9} };`

<code>arr[0][0] == 1</code>	<code>arr[0][1] == 2</code>	<code>arr[0][2] == 3</code>
<code>arr[1][0] == 4</code>	<code>arr[1][1] == 5</code>	<code>arr[1][2] == 6</code>
<code>arr[2][0] == 7</code>	<code>arr[2][1] == 8</code>	<code>arr[2][2] == 9</code>

- Logic Error: Don't use `arr[x, y]`
 - Wrong notation and ordering
 - Still compiles, evaluates `x,y` using the `,` operator, so really just `arr[x]`
- Logic Error: Note row-major ordering
 - So `arr[y][x]`

Jumping Around 2D Arrays

- All multi-dim arrays are really stored as one big contiguous array
 - Big O access time $O(1)$
- `int arr[3][3] = { {1,2,3}, {4,5,6}, {7,8,9} };`
- `int arr[3][3] = { 1,2,3,4,5,6,7,8,9 };`
- `int arr[3][3] = { 1,2,3,4,5,6,7,8,9 };`
- To get to the first element of the third row, its
 - $(\text{targetRow}-1) * \text{num_column_items} * \text{item_size_in_bytes};$
 - `arrayBasePtr + 2*(3)*4`
 - If the array starts at address 0 (ie, item "1" is at 0, item "2" is at 4)
 - The first item of the third row is at offset (24)

Addr (10)	Val
0	1
4	2
8	3
12	4
16	5
20	6
24	7
28	8
32	9

Passing Multi-Dimensional Arrays

- The second, third, ... `n` dimensions of arrays need to be specified to the function call
- All multi-dim arrays are really stored as one long contiguous array
- `int arr[3][3] = { {1,2,3}, {4,5,6}, {7,8,9} };`
- `int arr[3][3] = { 1,2,3,4,5,6,7,8,9 };`
- To get to the first element of the second row, we must know the number of columns:
 - Note zero-based counting for `targetRow` and `targetCol`
 - `arrayBase = ((targetRow * num_columns + targetCol) // "column_data_size" to get bytes`
 - `* (arr + ((x*totalColumns + 0)))` // here moves in strides of 4 due to int type
 - `arr[(1*3+0)]` // for the compiler to determine
 - `arr[y][x]` // convenient notation, yes?
- //to figure out how "wide" each row is, we need the col count
- `void printArr( const int data[][totalColumns] ) //need 2nd dimension`
 - `for(int x = 0; x < rows; x++) {`
 - `for(int y = 0; y < cols; y++) {`
 - `cout << data[x][y] << " ";`
 - `}`

gdb-g ↔ g++-g -c

valgrind --tool=memcheck --leak-check=yes

$$\begin{aligned}
 3y' + ty &= 4 \quad P(t) = \frac{t}{3} \\
 y' + \frac{t}{3}y &= \frac{4}{3} \quad u(t) = e^{\int \frac{t}{3} dt} = e^{\frac{t^2}{6}} \\
 e^{\frac{t^2}{6}} y' + \frac{t}{3} e^{\frac{t^2}{6}} y &= \frac{4}{3} e^{\frac{t^2}{6}} \quad u(t) = e^{\frac{t^2}{6}} \\
 \frac{d}{dt} \left[e^{\frac{t^2}{6}} y \right] &= \frac{4}{3} e^{\frac{t^2}{6}} \quad \int \frac{4}{3} e^{\frac{t^2}{6}} dt \rightarrow y \approx \frac{4}{3} e^{\frac{t^2}{6}} \int_0^t e^{\frac{s^2}{6}} ds
 \end{aligned}$$

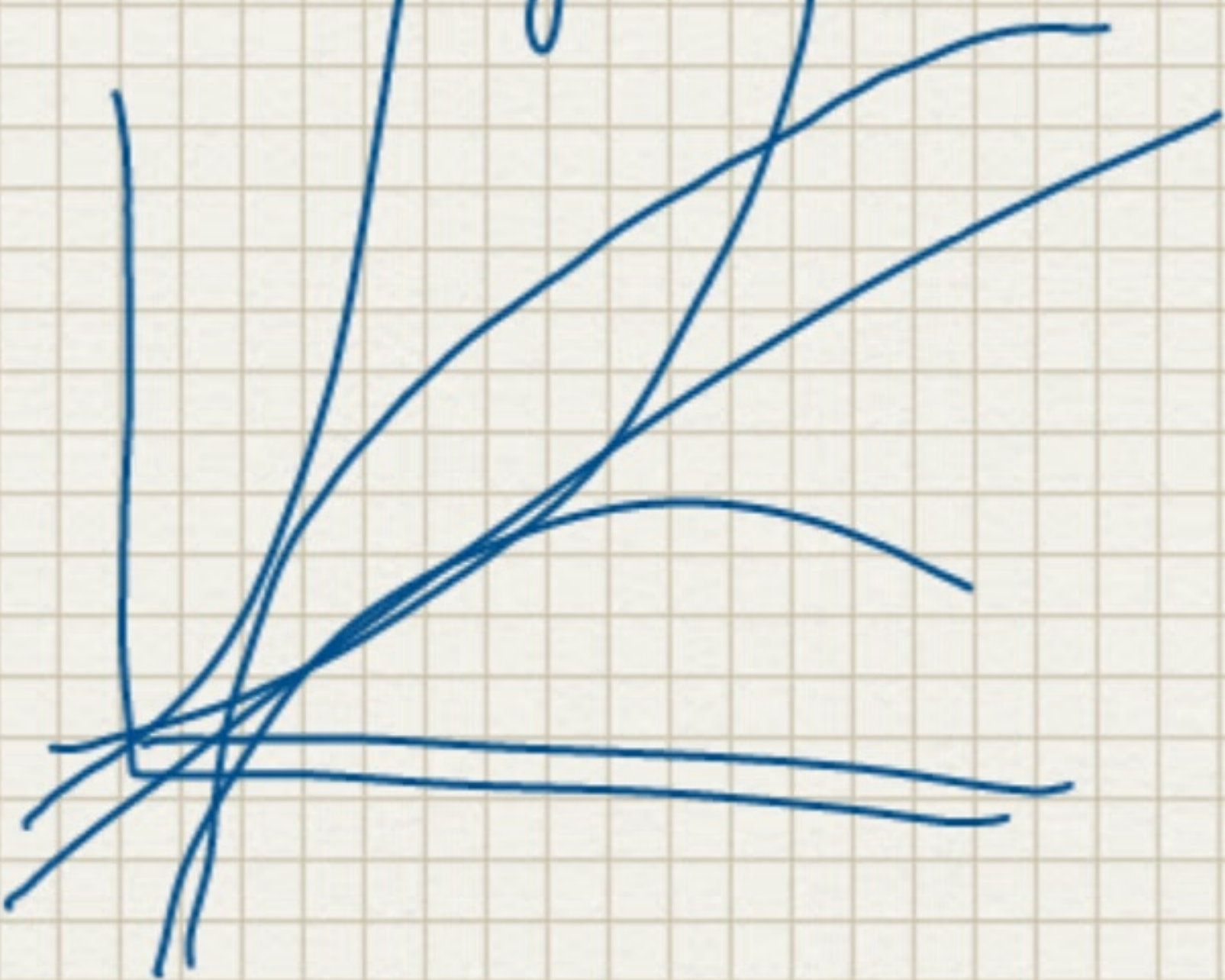
LOL DFC

Look @ Nash's Stack main.cpp Chapter 20 in Stroustrup
kill destructor, copy, assign

Could you help me
write the employee
destructor for me?
Employee: ~Employee() {

while n! = 0

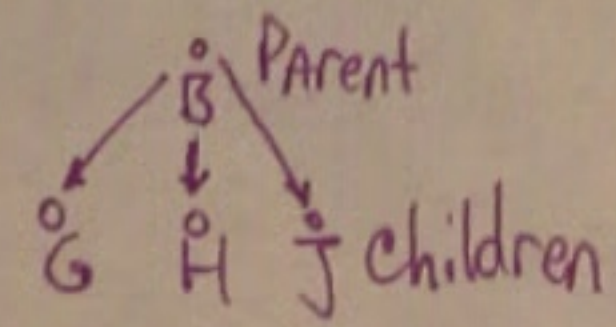
$$O(1) < O(\log n) < O(n) < O(n \log n) < O(n^2) < O(n^3) < O(2^n) < O(n!)$$



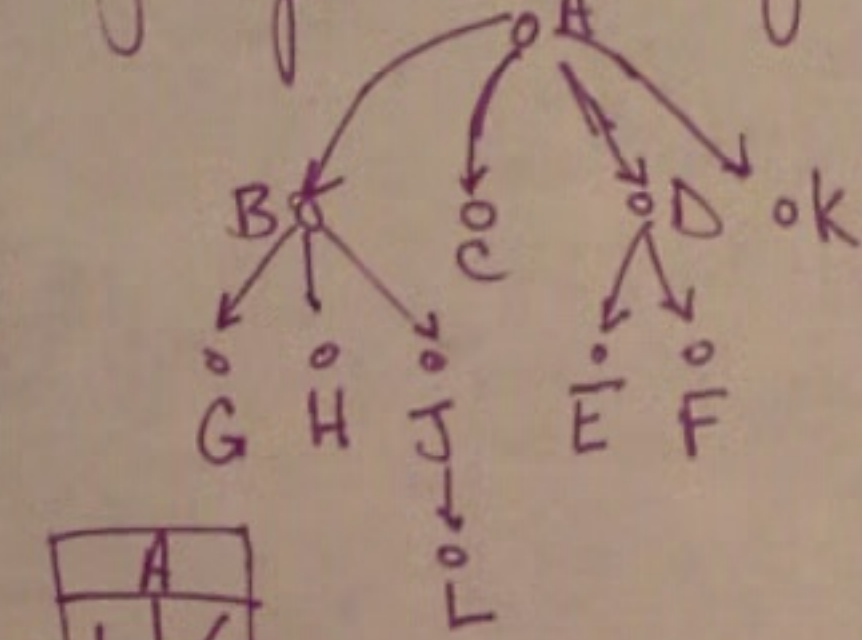
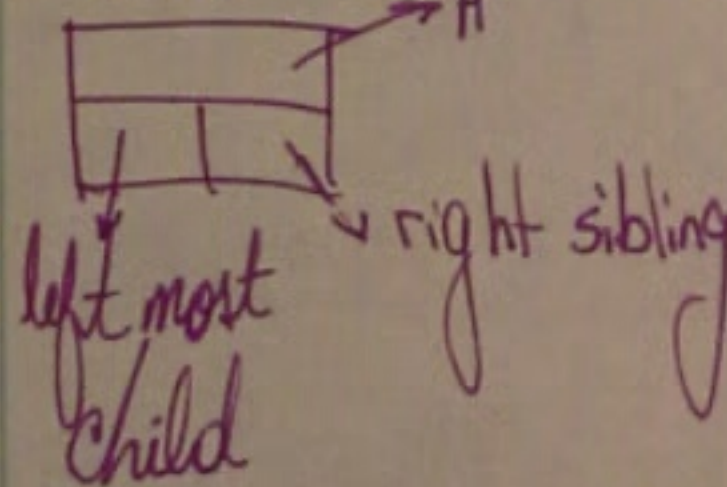
Disclaim unreliable methods.

○ Bubblesort for comment how-to. What code is doing $\Rightarrow$ not how.
end column comments $\rightarrow$ keep it short

Trees: A single node is a tree.
G H J



Common implementation of a general tree - leftmost child, right sibling

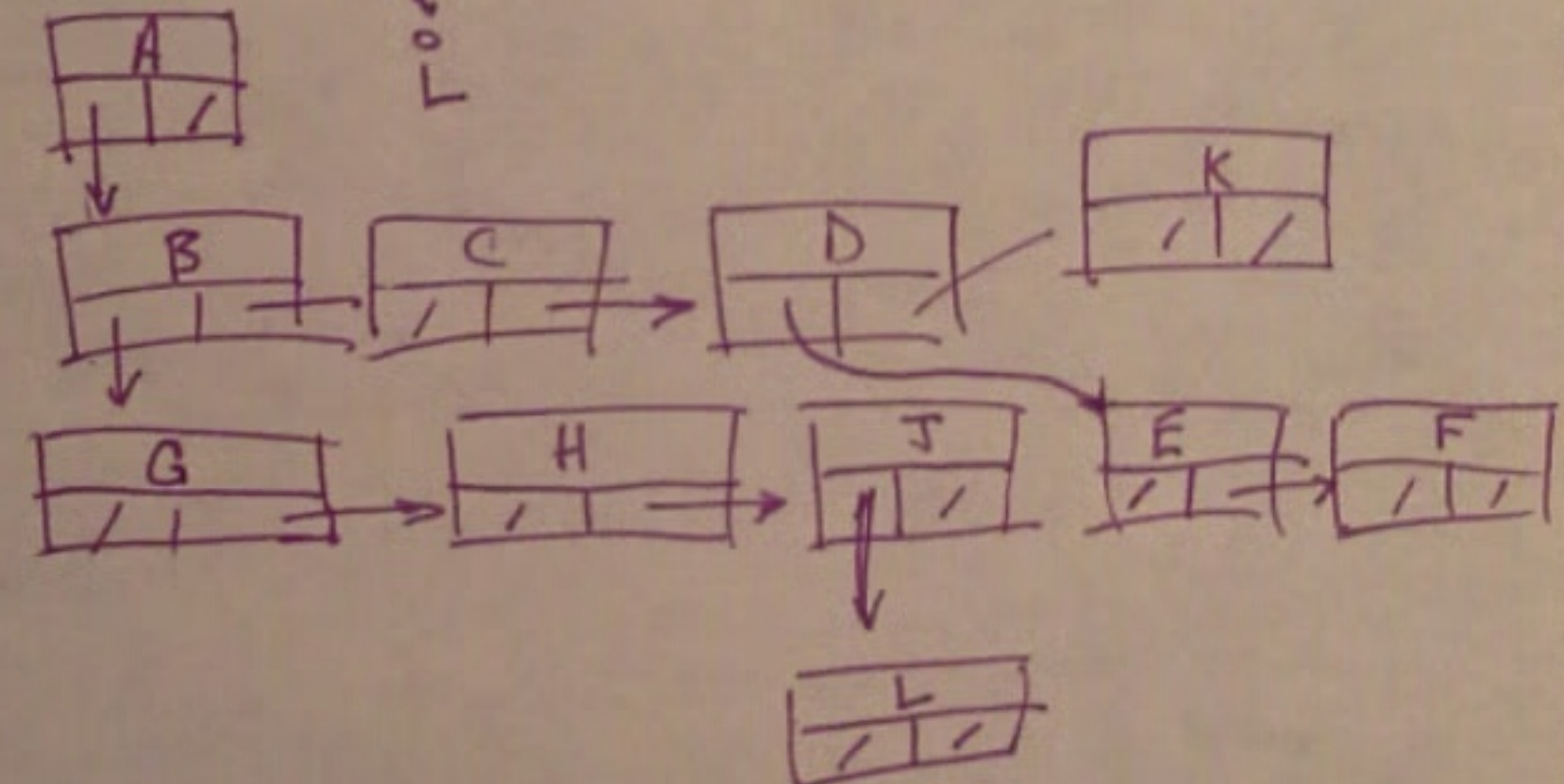


□ Expression Tree

□ BSP Tree

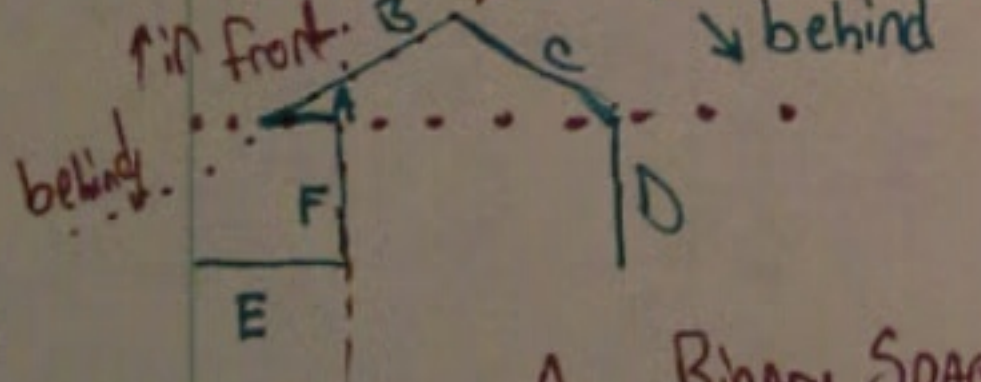
□ Huffman encoding algorithm

□ Binary Search Tree

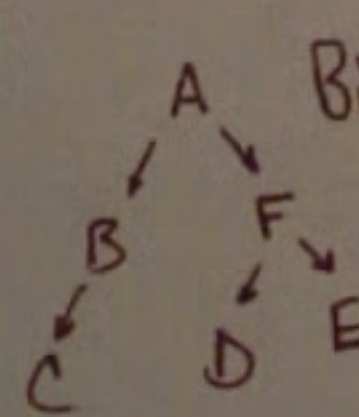


Encoding - Find a uniquely encoded and decoded code s.t. chars w/ higher freq have shorter encodings

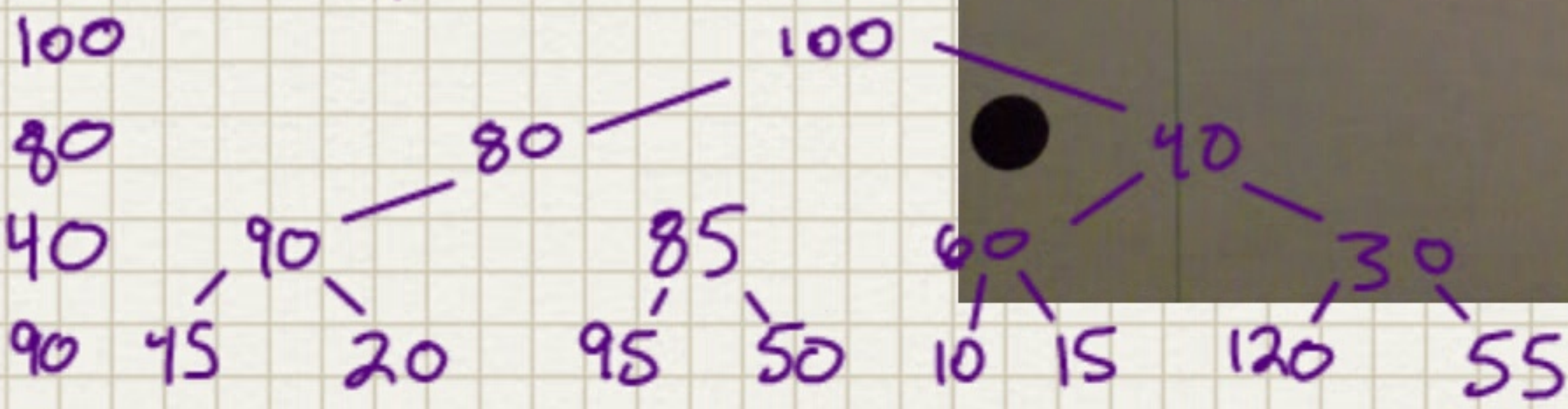
BSP Tree



Binary Space Partition



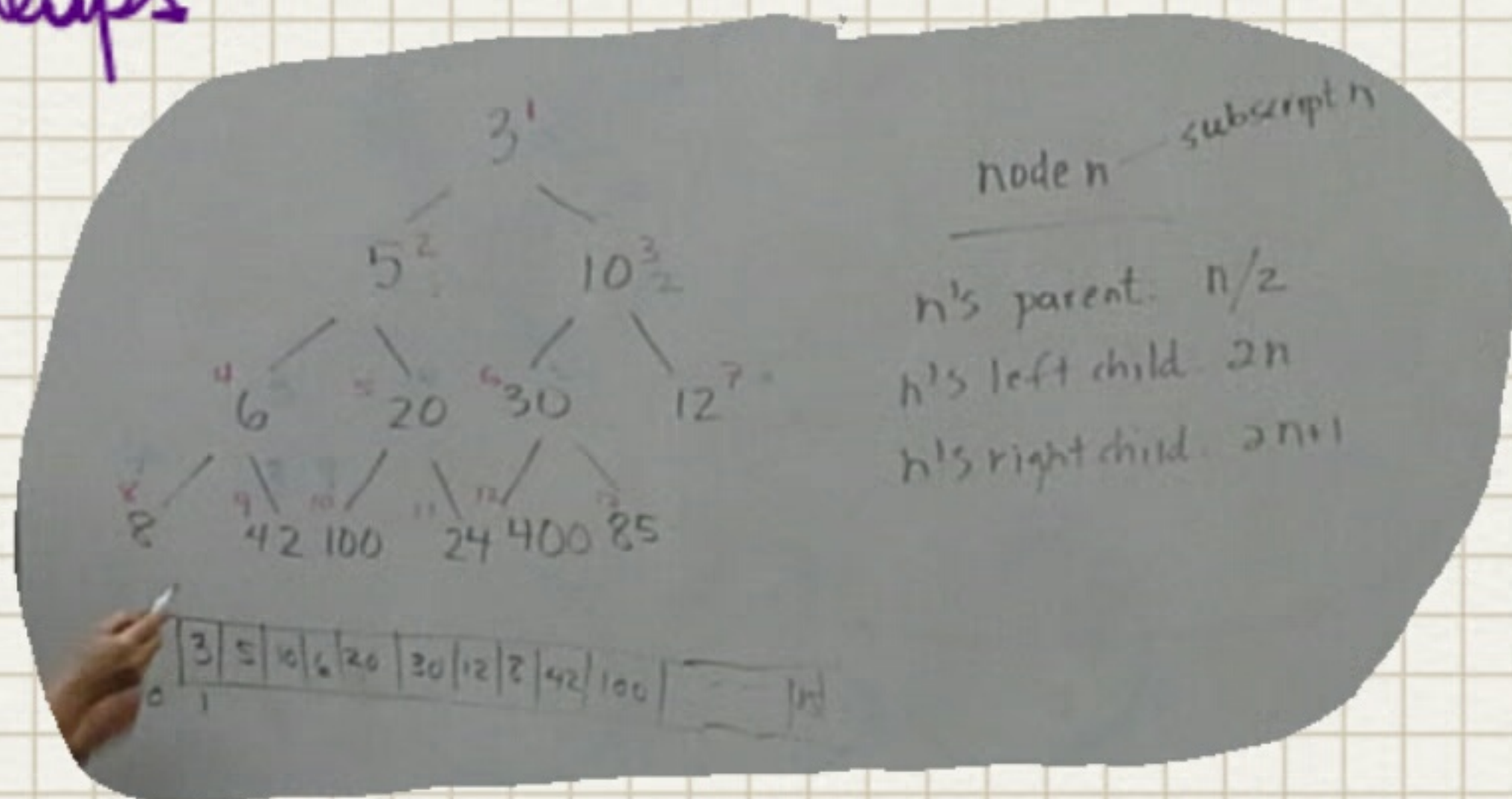
Build heap in linear time
must have data up-front
data - no particular order



sort / heapify lowest level

- Compare children, swap lowest
- little trees are heaps
- go up a level

sort next level up



node n subscript n

n's parent: $n/2$

n's left child: $2n$

n's right child: $2n+1$

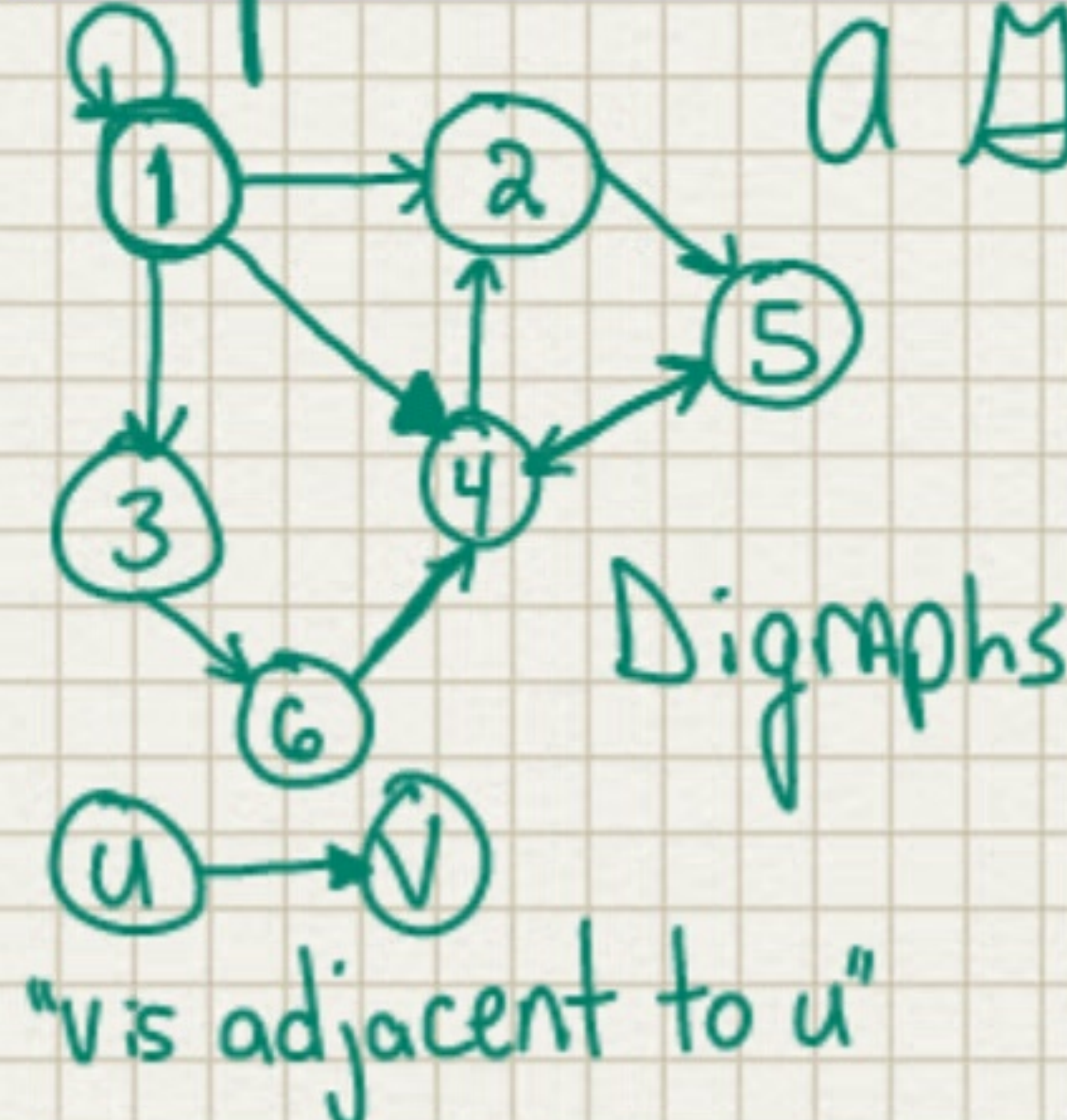
January 22, 2014

CS343 Lab 1 trends:

./a.out > file2

- comment all data members & functions
- over use of set/get
- output
- molest loop increment steps + mess w/ bounds.
- Linux → diff file1 file2
- Test getter/setters

Graphs



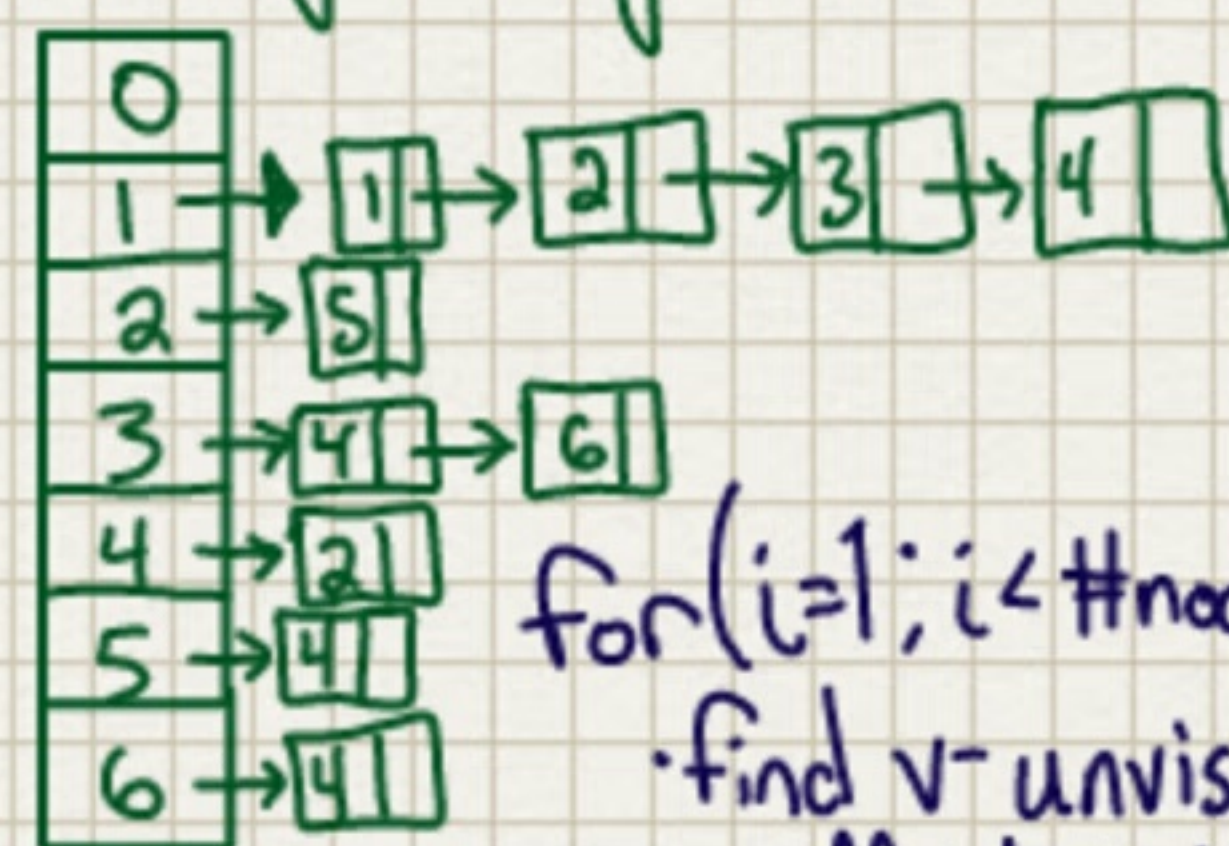
A Graph: (N, E) set of edges • edges connect vertices
 ↓
 set of nodes
 • edge is an ordered pair (u, v)

Implementation:

1. adjacency matrix = $A \rightarrow \# \text{ of Paths of length } 1$
 $\hookrightarrow A^2 \rightarrow \# \text{ of Paths of length } 2$
 $A^3 \rightarrow \# \text{ of " " " } 3$

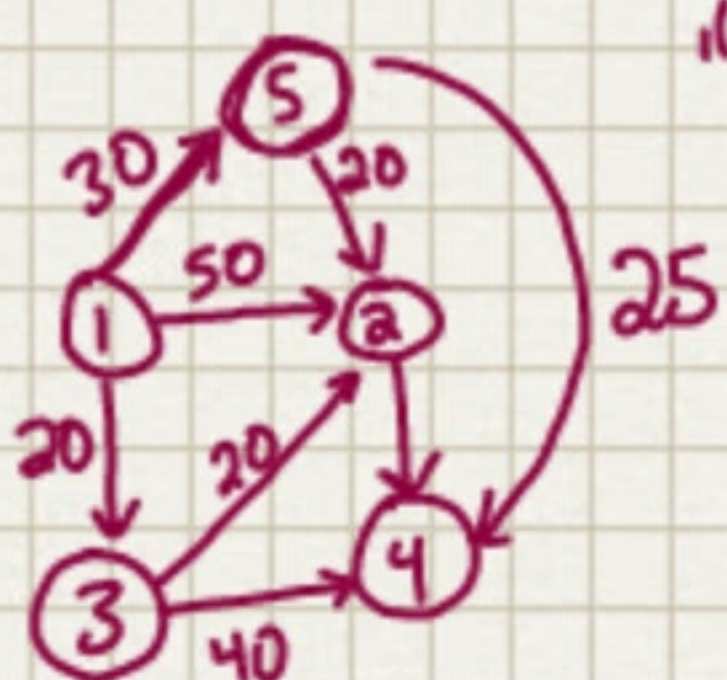
	1	2	3	4	5	6
1	1	1	1	1	0	0
2	0	0	0	0	1	0
3	0	0	0	1	0	1
4	0	1	0	0	1	0
5	0	0	0	1	0	0
6	0	0	0	1	0	0

2. Adjacency lists



Dijkstra's shortest Path

- finds the shortest distance from a source node to every other node.



• Given a source •

	1	2	3	4	5
1	-	50	20	-	30
2	-	-	-	10	-
3	-	20	-	40	-
4	-	-	-	-	-
5	-	20	-	25	-

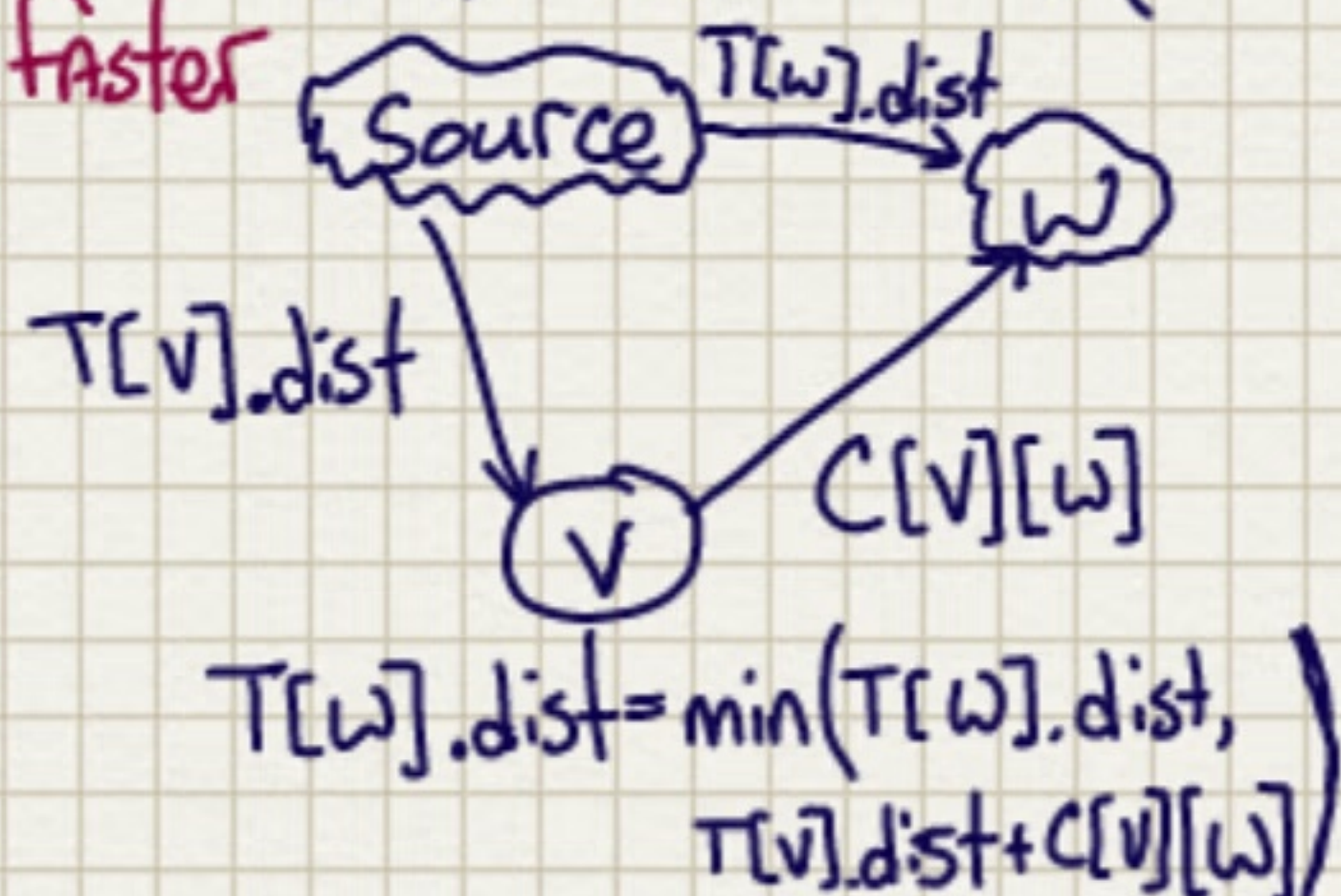
Cost Matrix

"Greedy Algorithm"

- Consider the source
- Consider lowest cost edges
- mark lowest edge vertex
- Try next-lowest-edge vertex
- Try to use other paths to try to get there faster

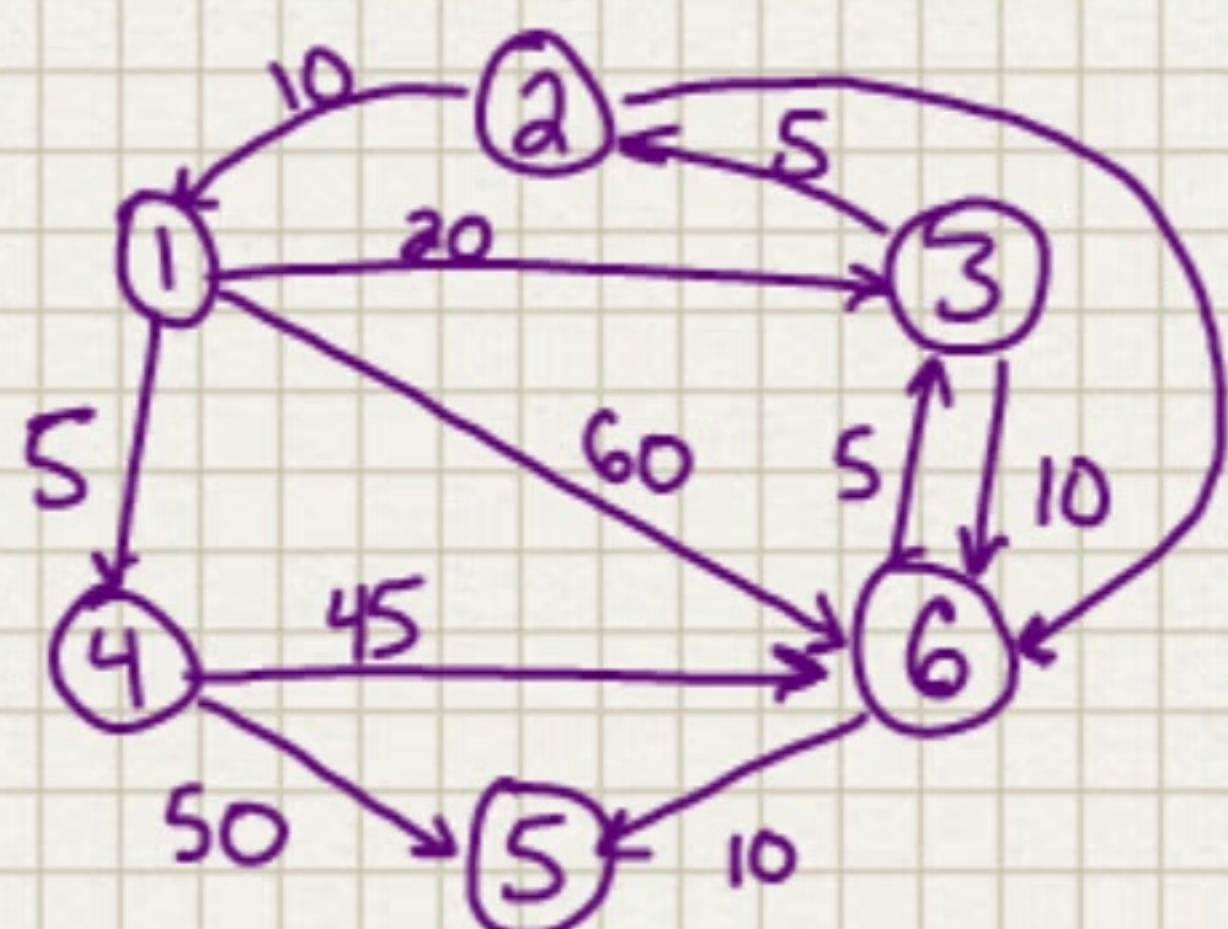
$T[w/e].dist$	1	2	3	4	5
	0	∞	∞	∞	∞

for $(i=1; i < \#nodes-1; i++)$ {
 • find v - unvisited with smallest path distance.
 • MARK v visited
 • for each w adjacent to v
 if (w is not visited) {



1/27/2014

hold a bool to validate your graph - bool says if graph has changed.
 ↳ ADT → a graph class → Dijkstra's just a method.



Source = 1

T[w].dist

	1	2	3	4	5	6
5	0	∞	∞	∞	∞	∞
V=1	∞	20	5	∞	60	
V=4	∞	20	55	50		
V=3	25		55	30		
V=2			55	30		
V=6				40	30	

T[w].Path

	1	2	3	4	5	6
	0	0	0	0	0	0
V=1	0	0	1	1	0	1
V=4	5	0	1	5	4	4
		3	5		4	3
					6	5
						3
	0	3	1	1	6	3

T[w].dist

↳ cost from source to w

	0	25	20	5	40	30
	1	2	3	4	5	6

Want path from 1 to 5

5

T[5].path = 6

T[6].path = 3

T[3].path = 1 → We can stop myeahh.

T[1].path = 0 → 1, 3, 6, 5

from to to

T[to].PATH = C

T[C].PATH = B

T[B].PATH = A

T[A].PATH = from

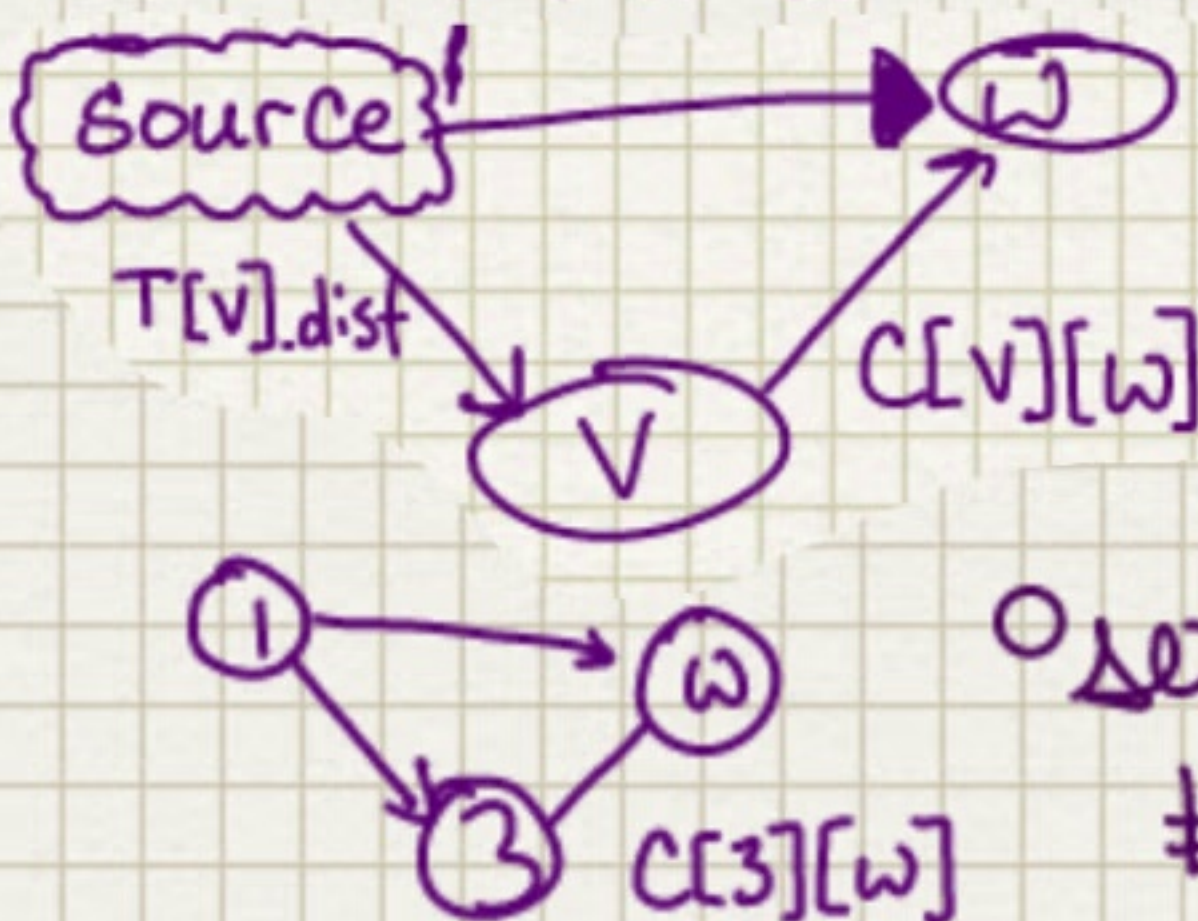
From ABC To

int i = to;

curr = T[from][to].PATH;

cout << curr

curr = T[from][curr].PATH



set width

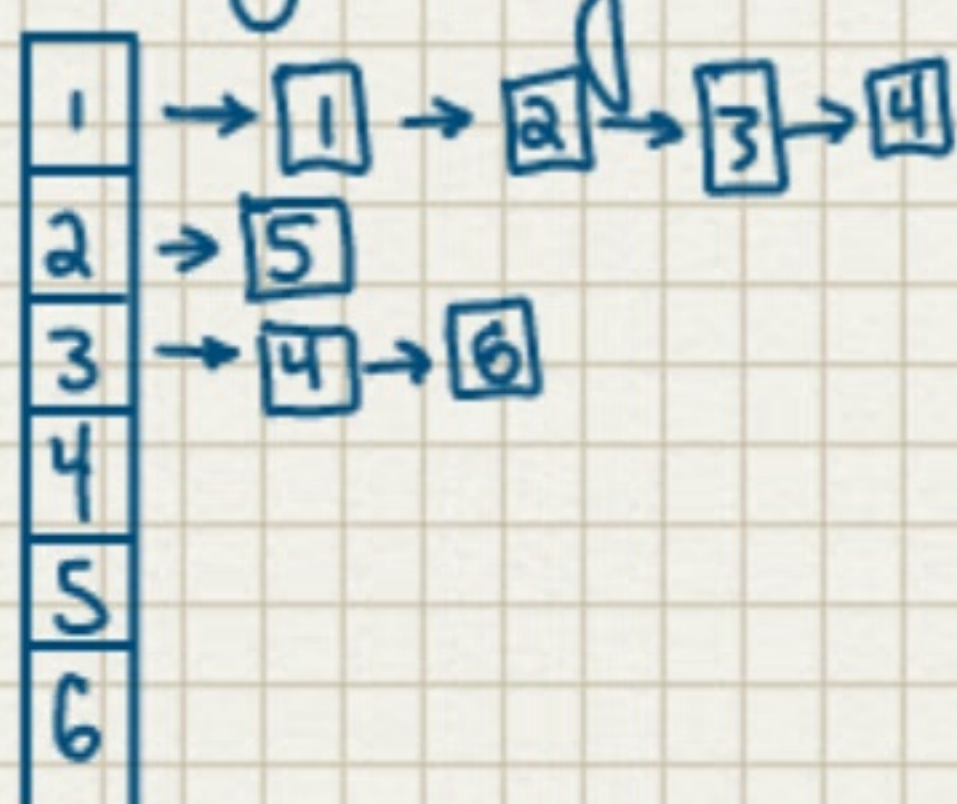
#include <iomanip>

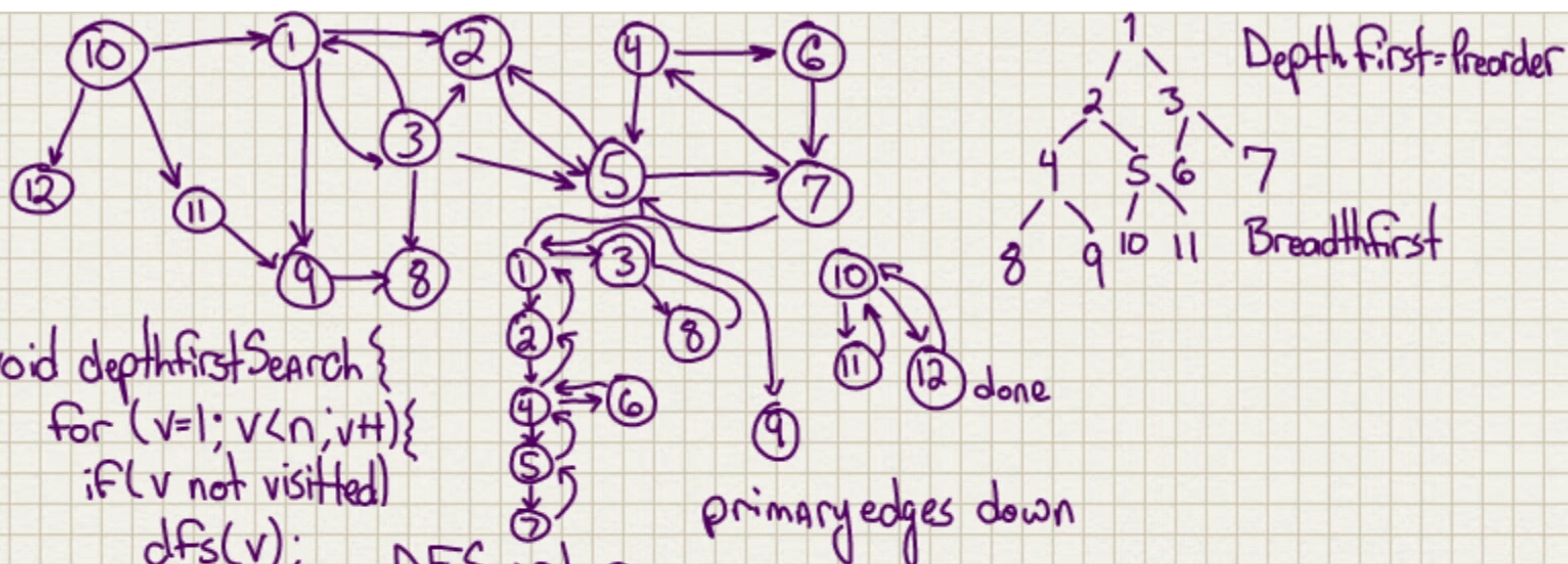
cout << setw(5) << 10

----- 10

----- 100

2. Adjacency List



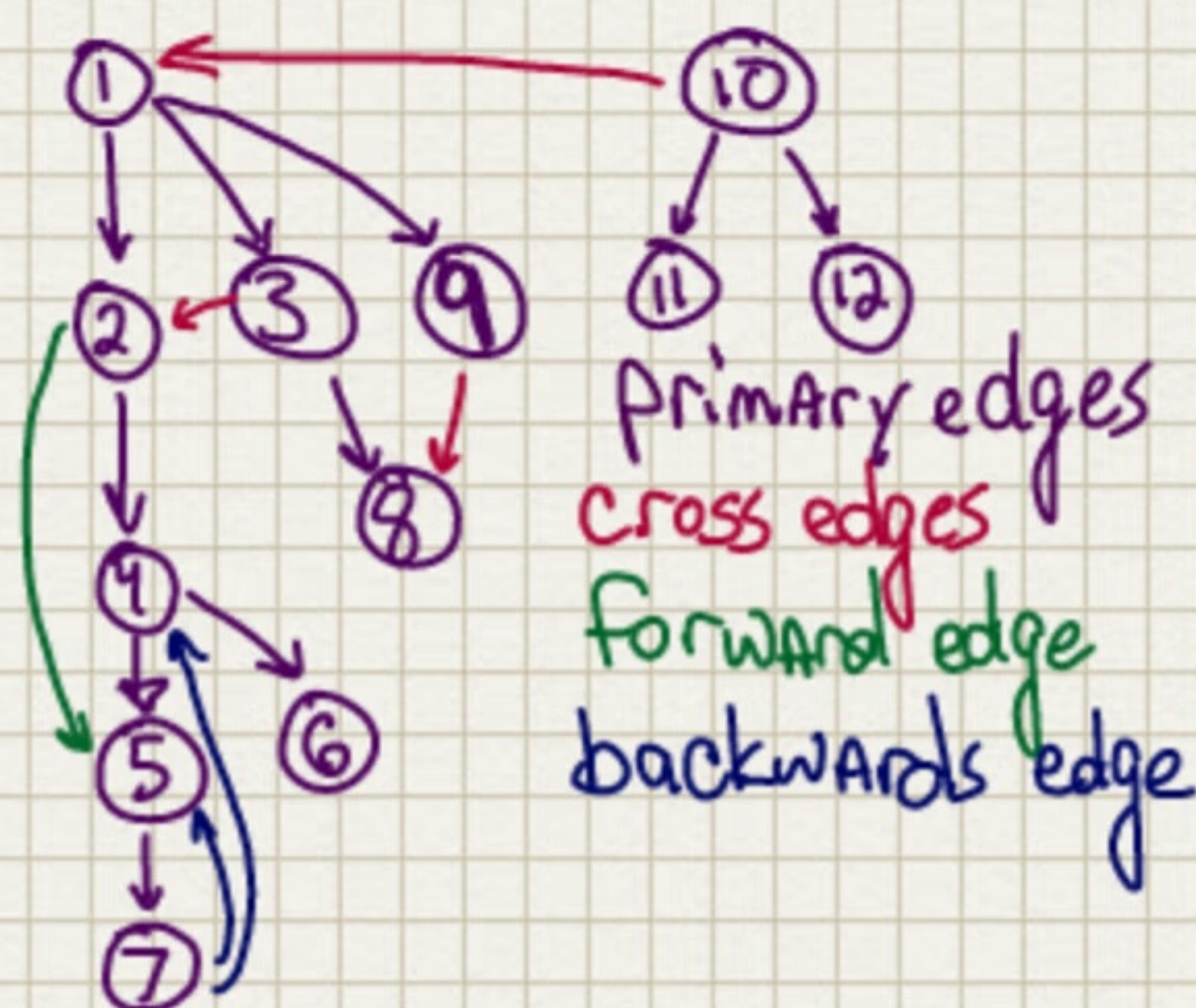


```
void depthfirstSearch {
  for (v=1; v<n; v++) {
    if (v not visited)
      dfs(v);
  }
}
```

```
void dfs(v)
  mark v visited.
  do whatever with v
  for each w adjacent to v
    if (w not visited)
      dfs(w)
```

DFS order

1 2 4 5 7 6 3 8 9 10 11 12



Breadth First Search

```
void bfs(v)
  mark v visited
```

```
q.enqueue(v)
```

```
while q not empty
```

```
  x = q.dequeue();
```

```
  do whatever with x.
```

```
  for each w adjacent to x.
```

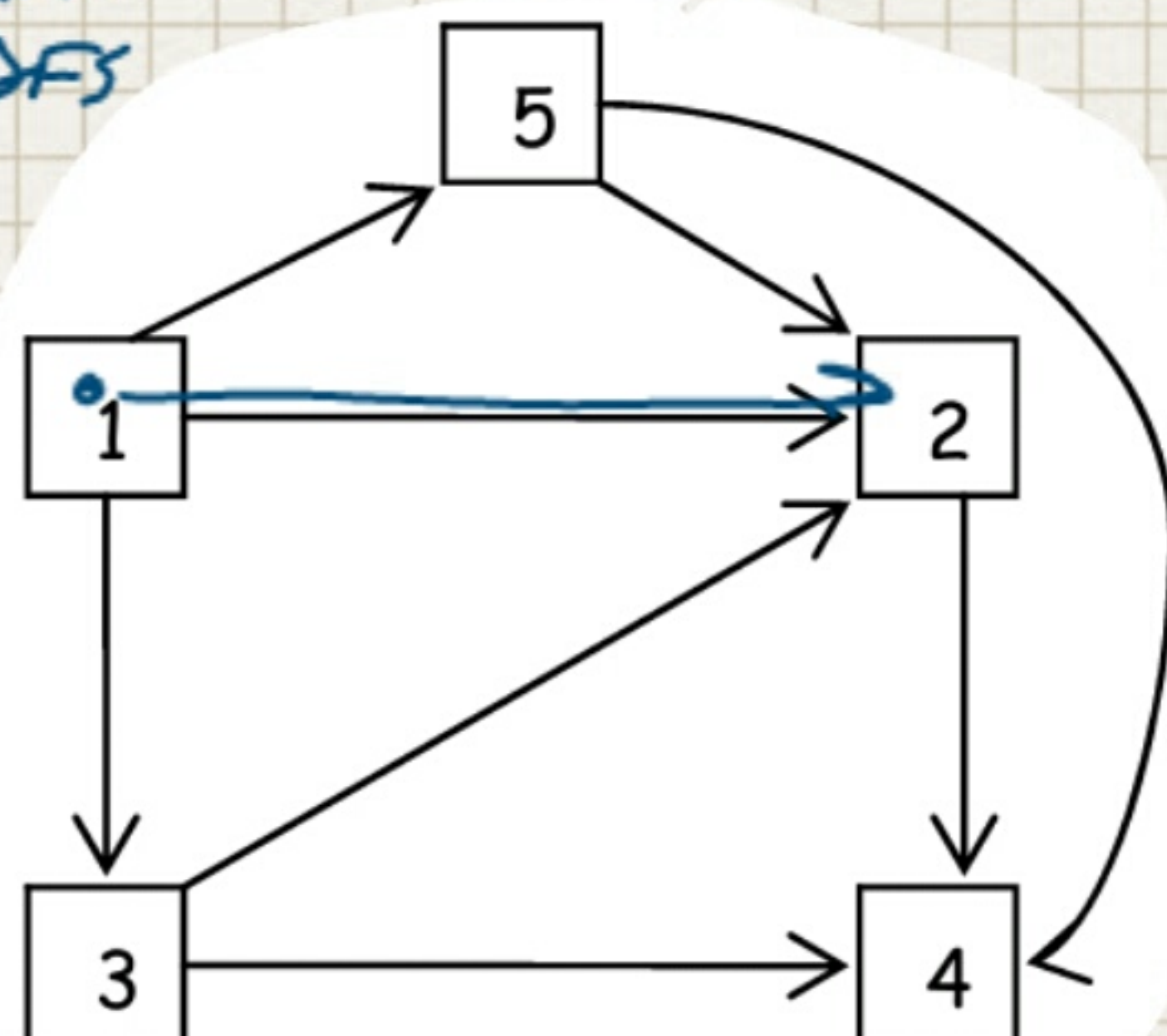
```
    if (w not visited)
```

```
      mark w visited
```

```
      q.enqueue(w)
```

start at first node defined $\rightarrow 1$
 $\rightarrow$ last edge of this node
 $\rightarrow$ mark
 $\rightarrow$ DFS

5
 Aurora and 85th
 Green Lake Starbucks
 Woodland Park Zoo
 Troll under bridge
 PCC
 1 5
 1 3
 1 2
 2 4
 3 4
 3 2
 5 4
 5 2
 0 0



unsigned problem →
constructor →
zero poly → not printing zero
Read!

Test cases
- negative exponents
- negative coeffs

Don't overcomplicate → follow
Zander's more closely

1/29/2014

AVL Trees → Definition of Balanced: $\forall$ node $\in$ tree, the height of its left
and right subtree differs by at most one.

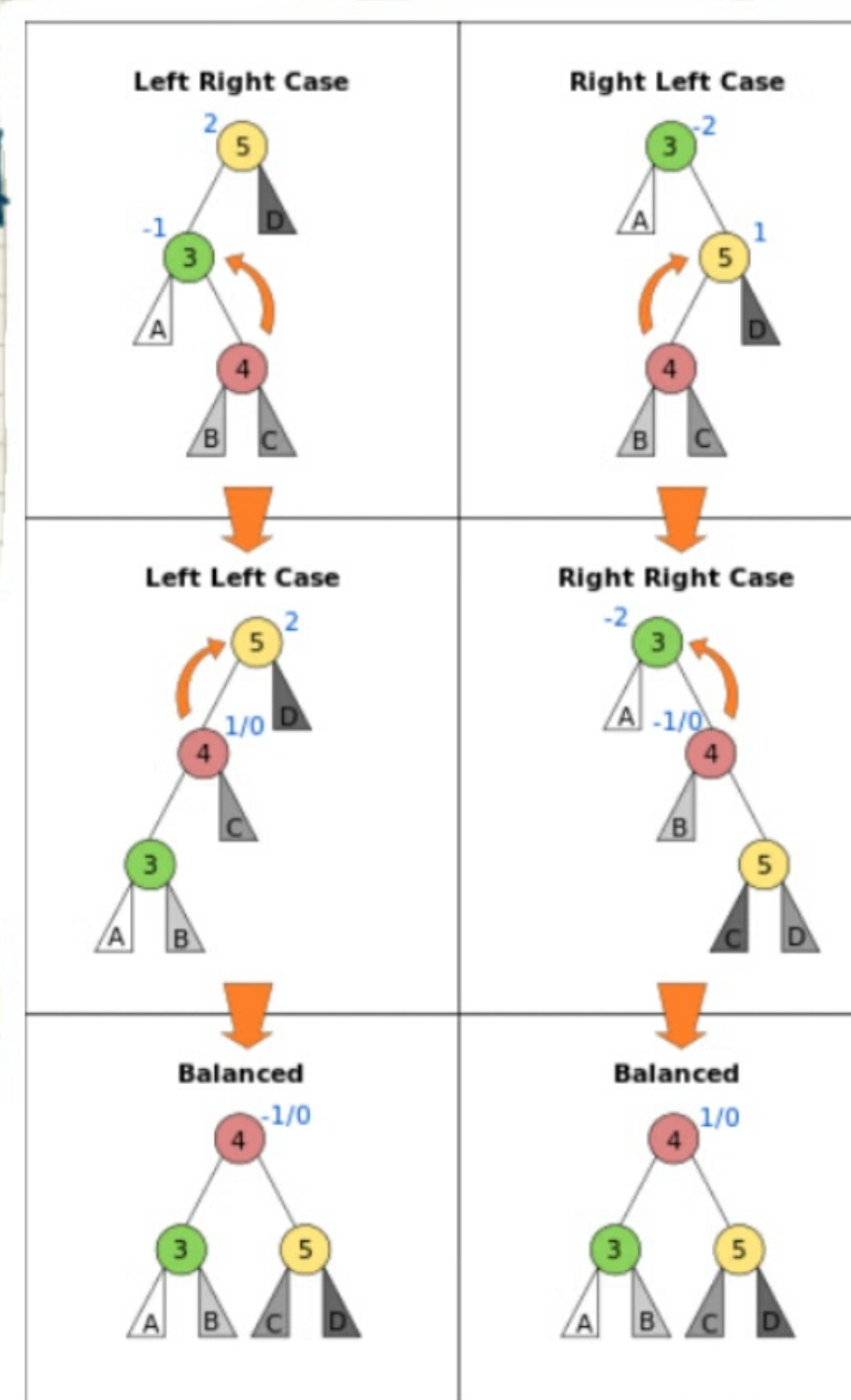
"Rotations": 4 kinds:

→ single rotation

- left left
- right right

→ double
rotations

→ left right
→ right left



http://en.m.wikipedia.org/wiki/AVL_tree

Huffman coding

Trees - B.S.P

- Huffman Coding

- Binary Search tree

10 points? 10 minutes

20 points? 20 minutes!

↳ Actually "take" practice / exam.

2-3 tree

↳ B-tree: Balanced trees

- the root is either a leaf or has between 2 and m children.

All non-leaf nodes have between $\lceil \frac{m}{2} \rceil$ and m children.

Poly class → know the code

↳ review Labs 1 & 2

Execution tree!!!

↳ recursion on paper

Heaps → Problem based

↳ short answer

↳ not too much code

↳ Data structures

Graphs → Dijkstra's

→ "Dijkstra doesn't work on negative edges"

Depth first

→ AVL insert "by hand"

Breadth first

Balance Trees ZOMG Complexities!

↳ AVL + 2/3 + "Tris"

↳ not too codey → insert → left-left

↳ Up to 1/3/2014

2-3 trees ^{insert is} Typically recursive.

All nodes are leaves or have 2 or 3 children.

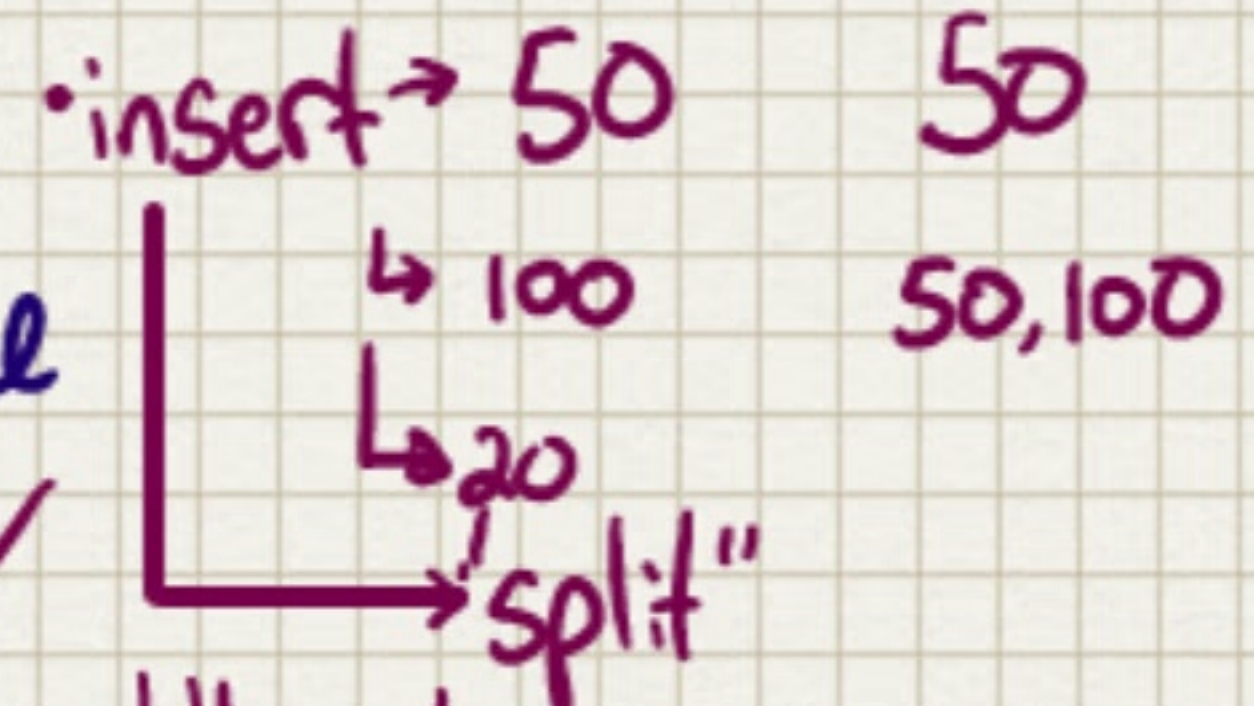
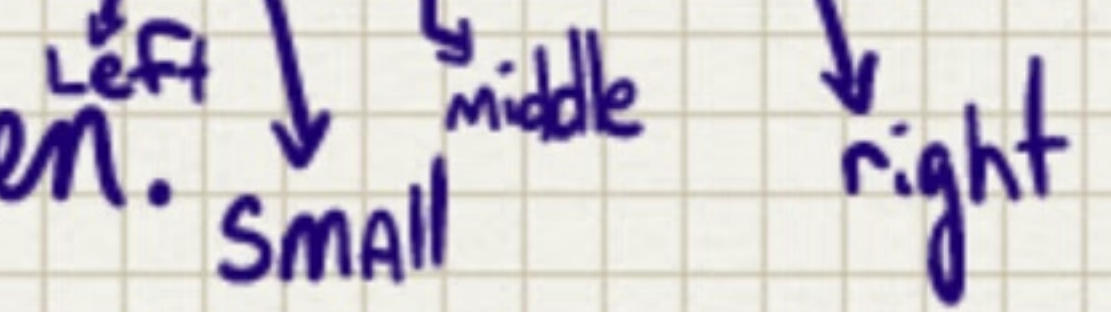
→ node has only 1 value → 2 children → right pointer = null.

→ node has 2 values → has 3 children.

Items that are smaller go in left subtree

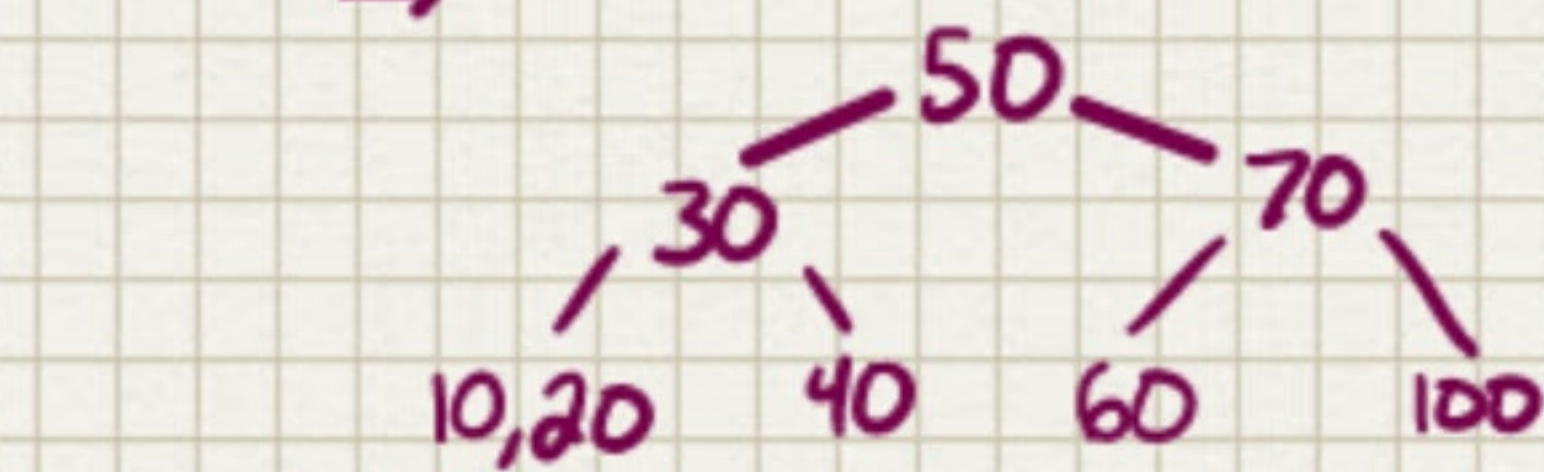
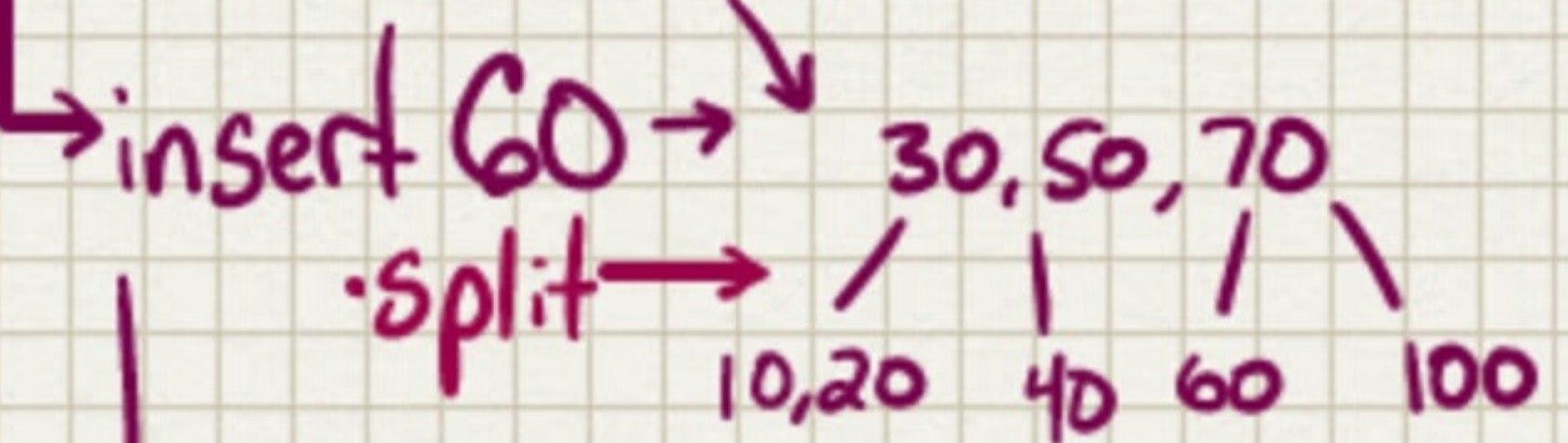
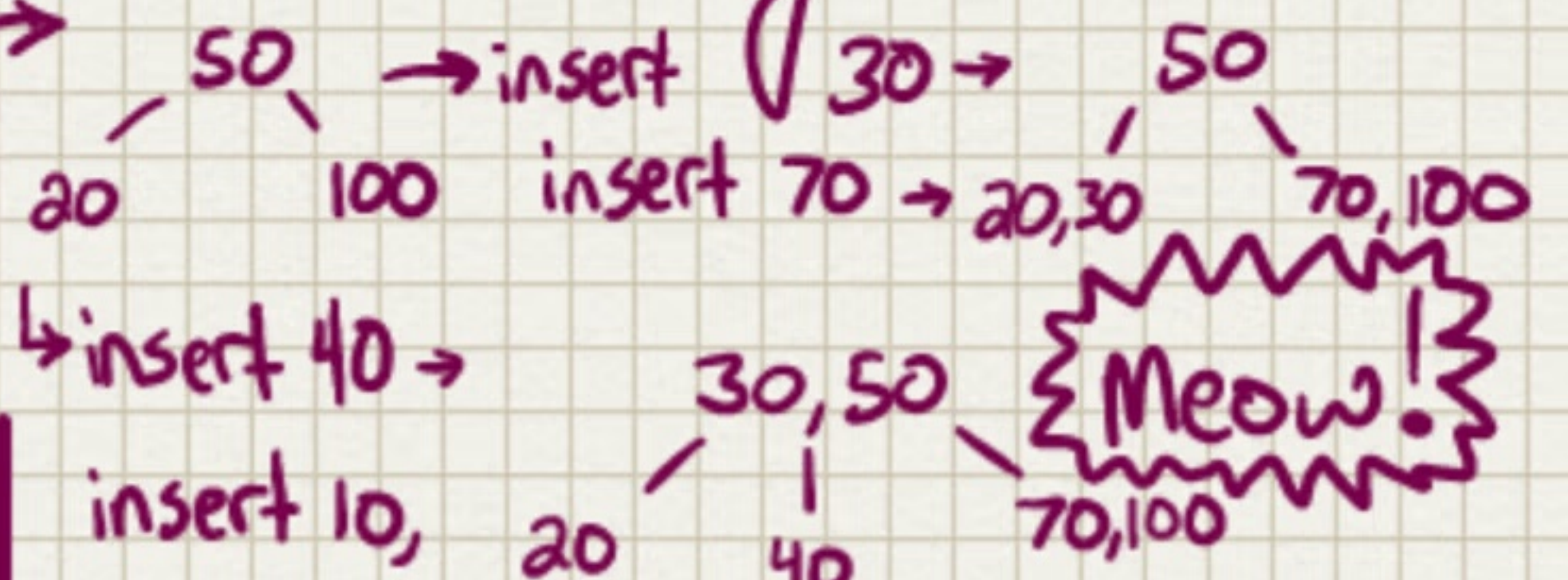
Items > S, < L go in middle subtree

Items > L go in right subtree.

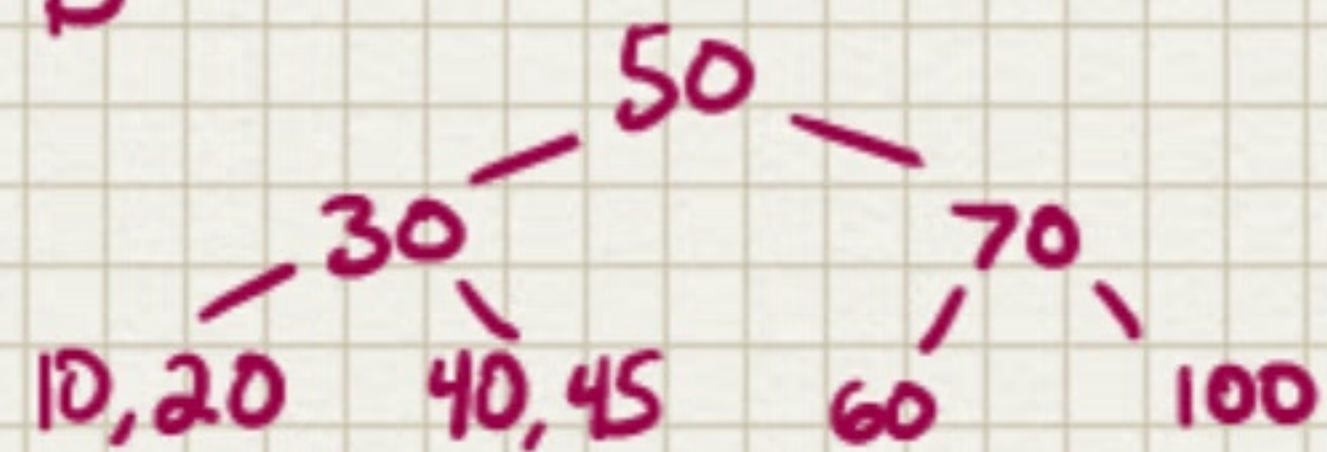


1. middle value "moves up"

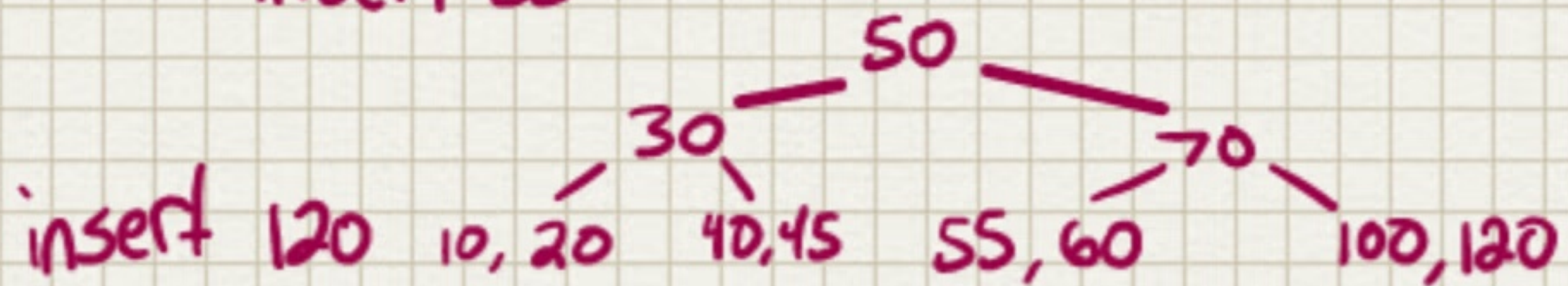
2. surrounding nodes become children



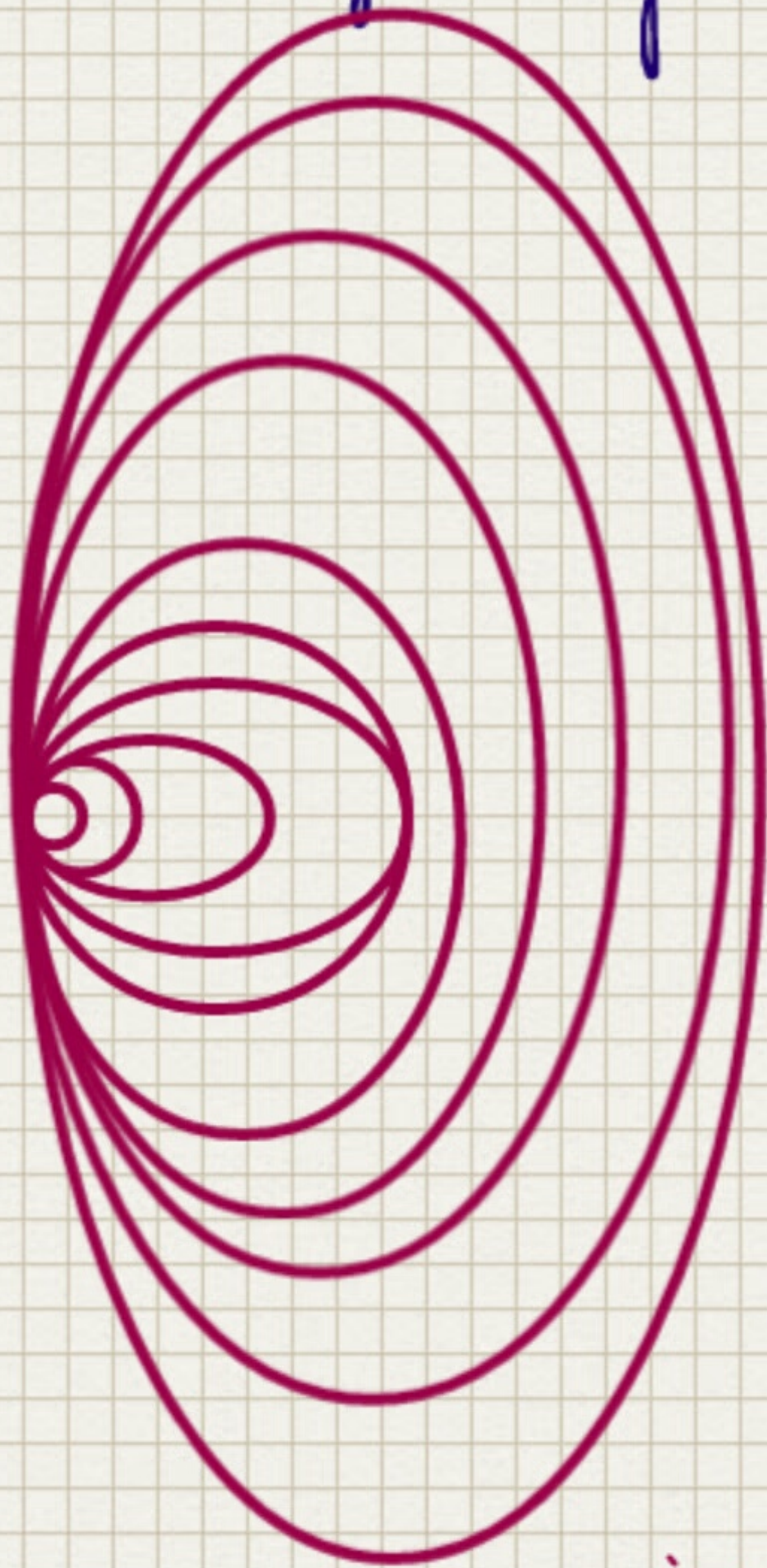
insert 45



insert 55



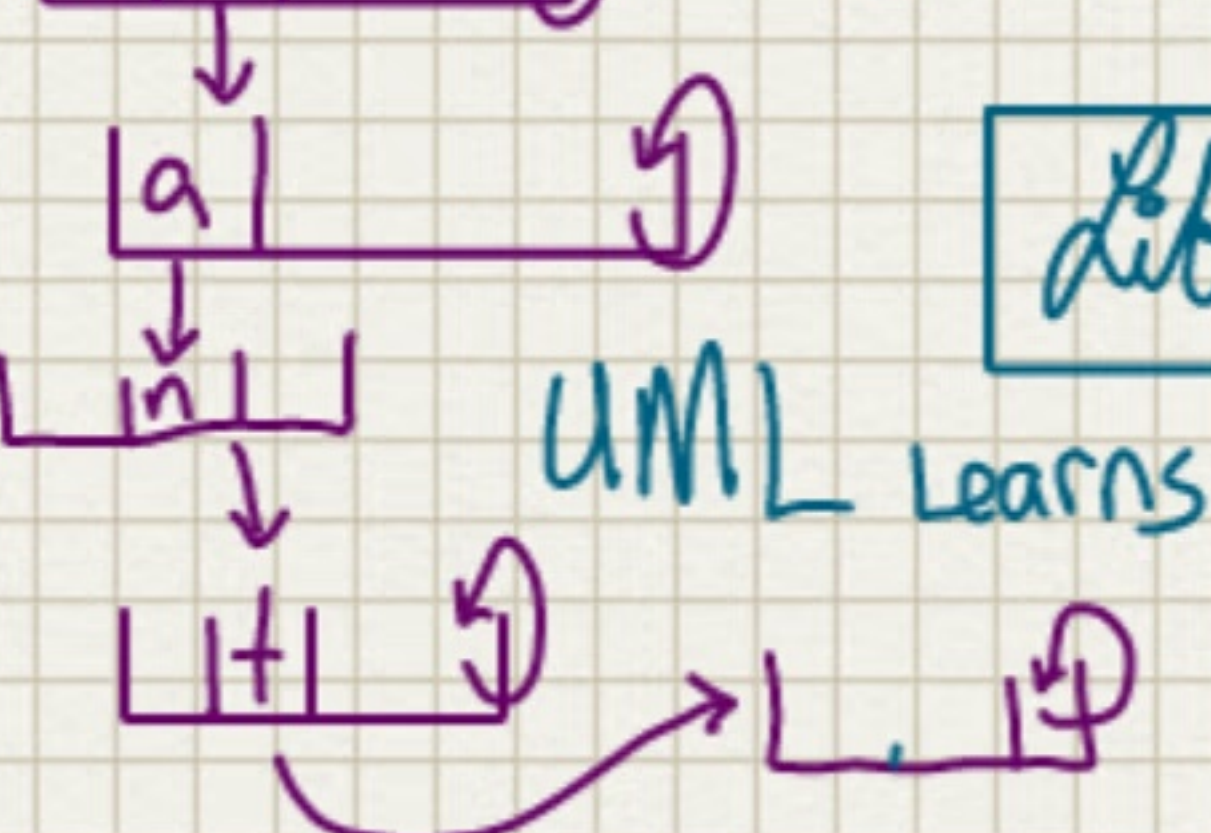
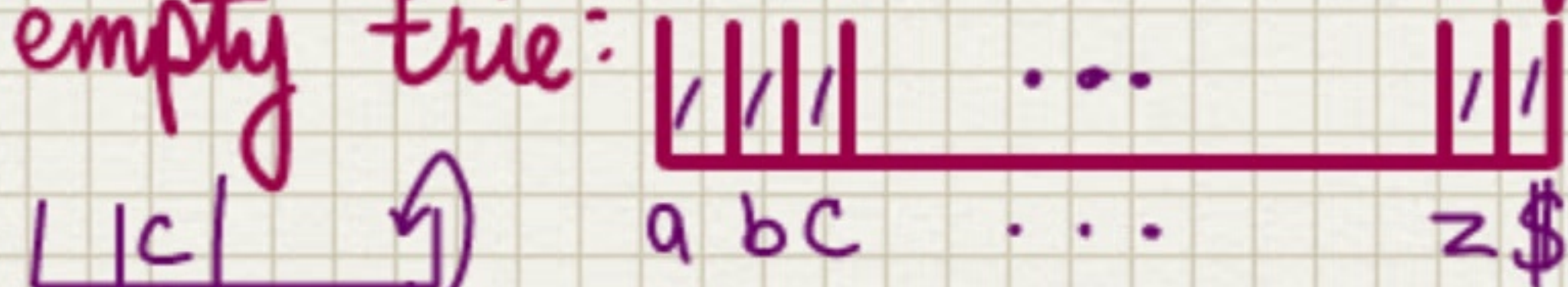
insert 65 →



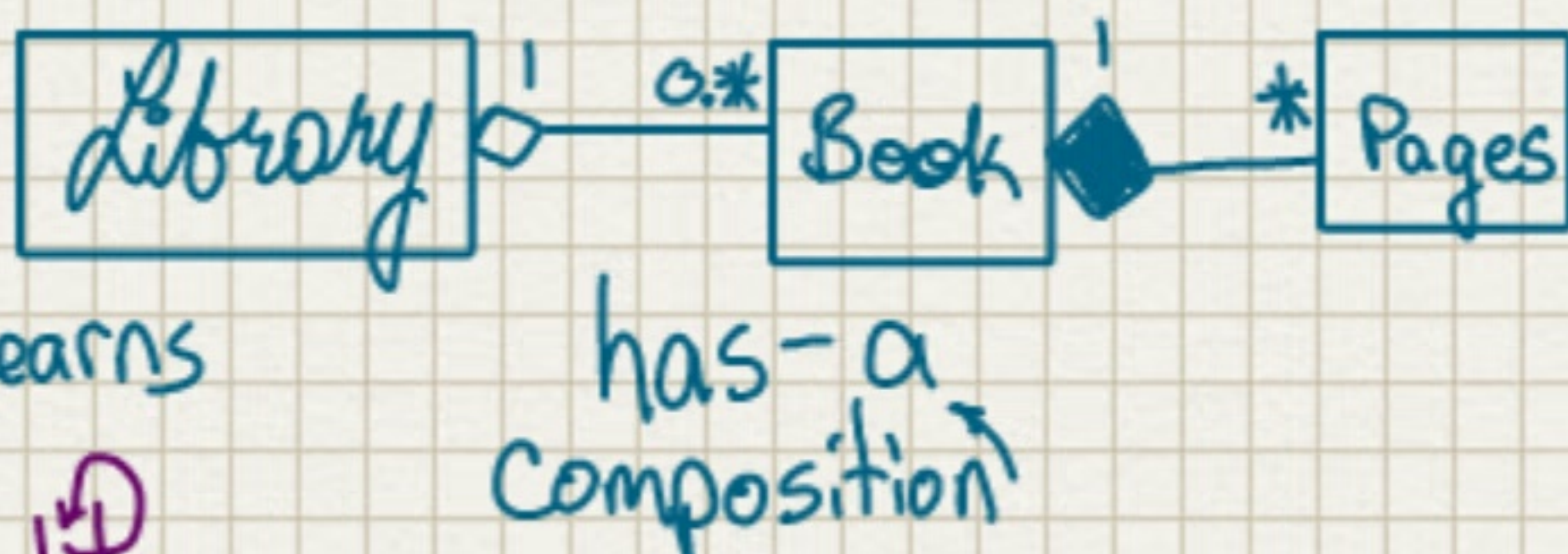
Trie (reTRIEval)

one way to implement a dictionary (or any n-bit string)

empty trie:

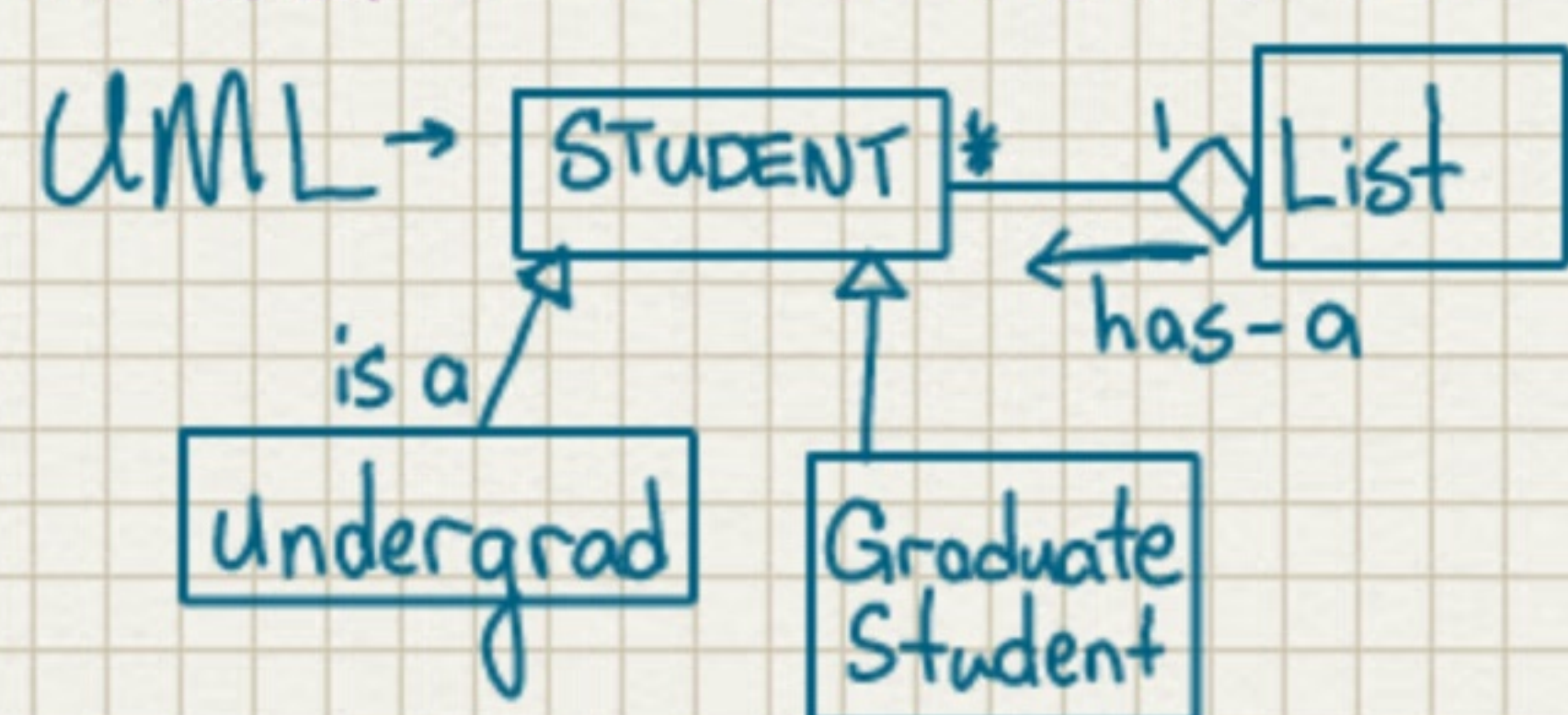


UML Learns



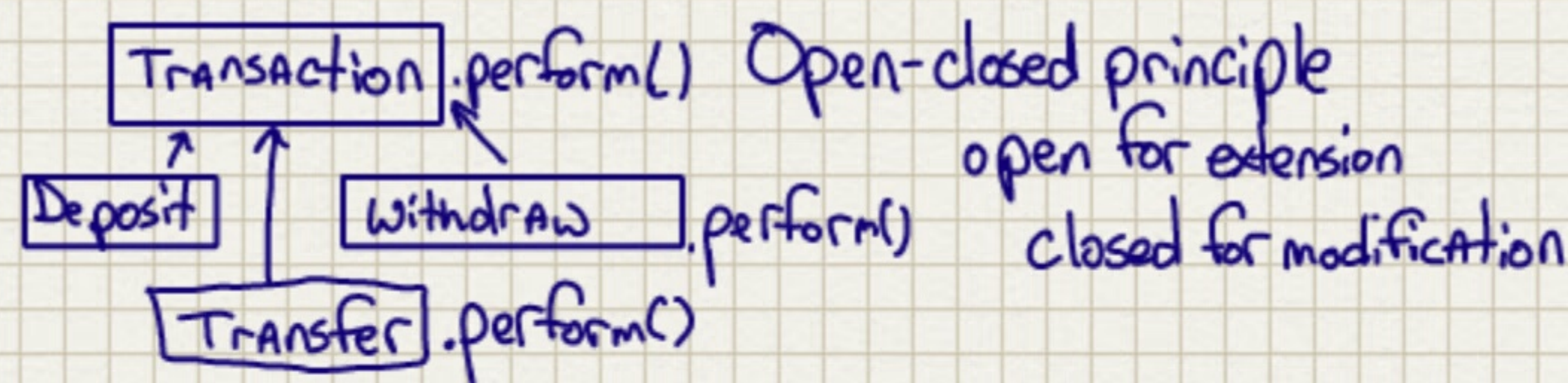
Insert for a trie → Hashing → if NULL new node* → update cur to pt to next.
ptr to array

Inheritance → NOT Different in C++



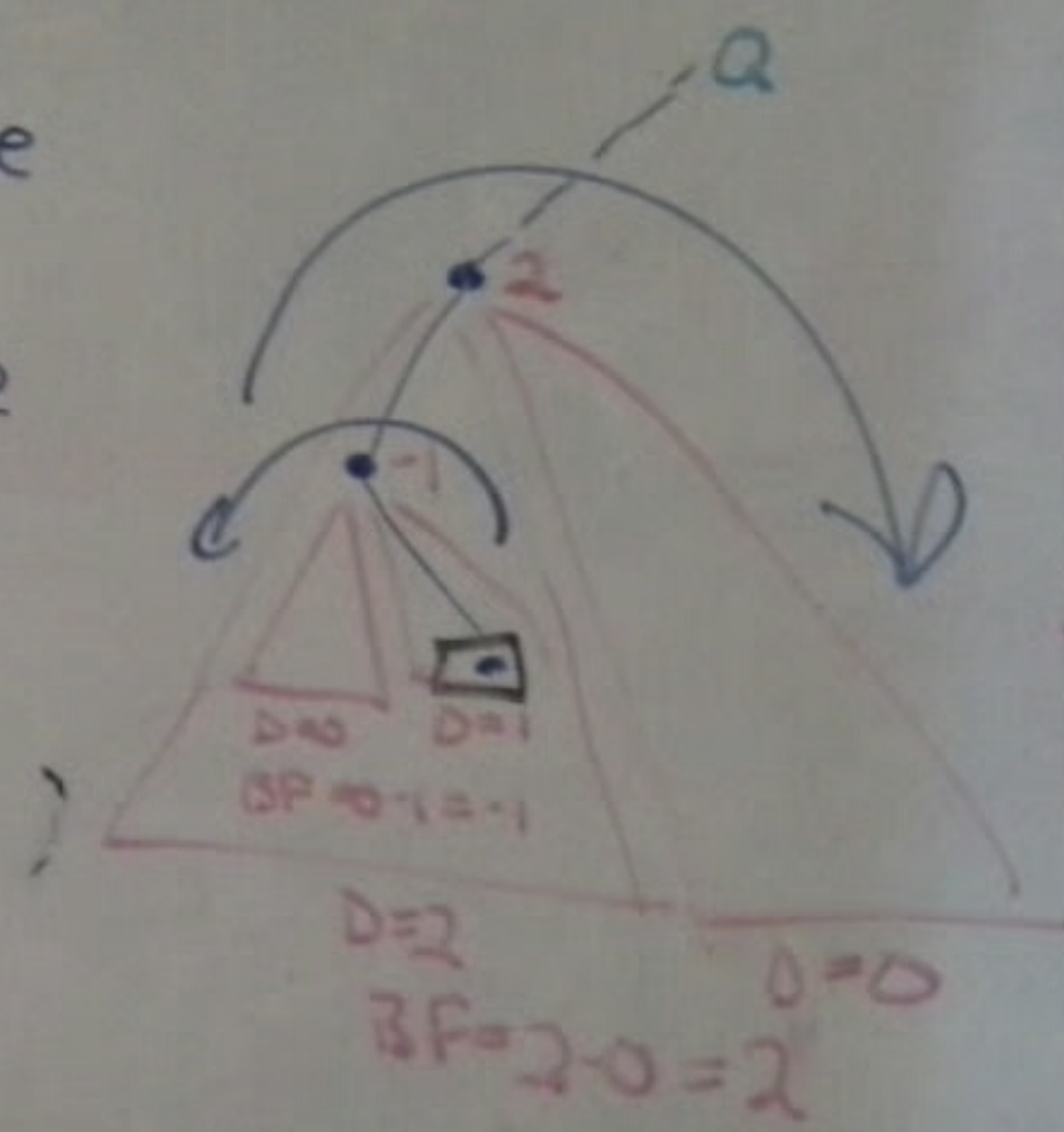
Aggregation/Composition

• base class ↔ super classes
↳ derived class ↳ subclasses



Trie Node -
 char letter;
 Trie Node *[] children; // size 27

$\begin{matrix} \text{RR} \\ \text{LL} \end{matrix} \} \text{Single}$
 $\begin{matrix} \text{LR} \\ \text{RL} \end{matrix} \} \text{double}$



Wed. 26 Mar

Call about

6 Feb @ 3:00
 Thursday

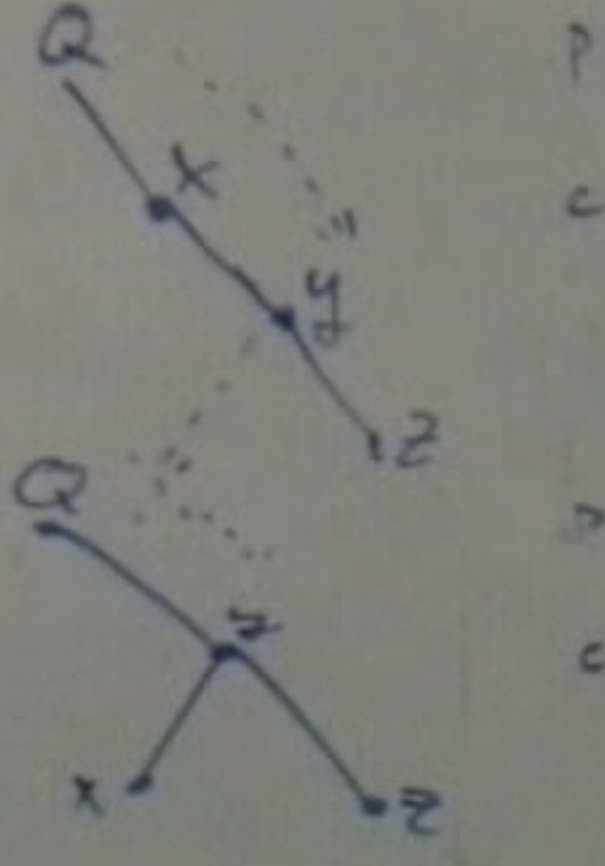
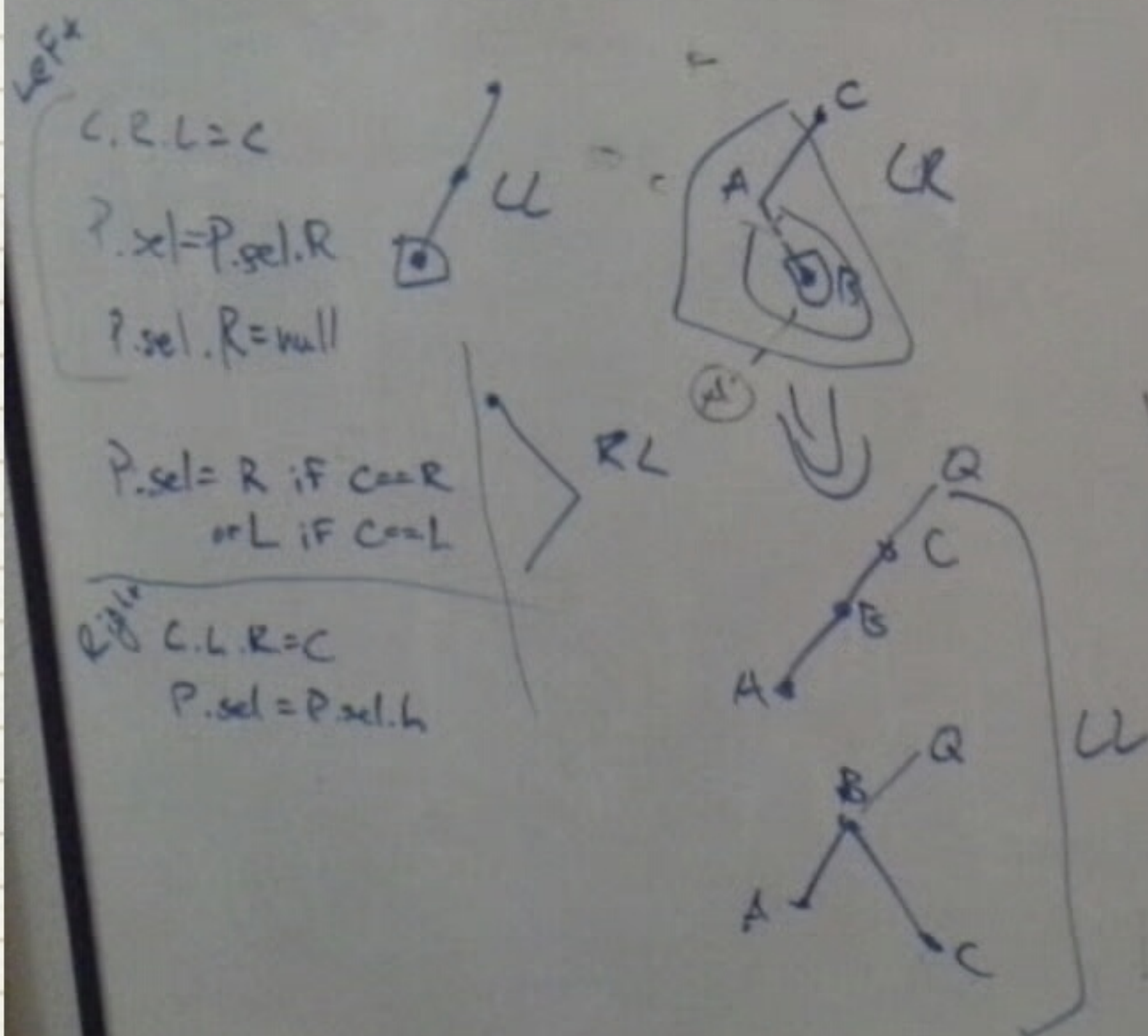
UW Seattle -

Med 206 520

206 744 311

arm x-ray

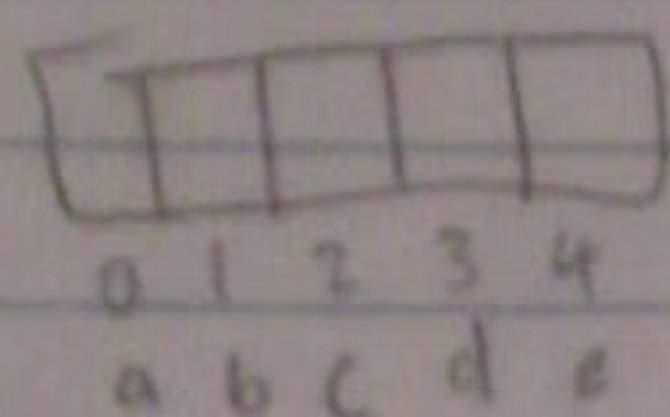
Phone Interview



hash table - fast lookup

2/5/14

Type: char letter [a]



int index = letter - 'a'; (4)

if NULL, create new??

else, current to next pointer??

Design-class diagram, simple

Put all in .h file for Lab 4 design turn-in.

UML Notes

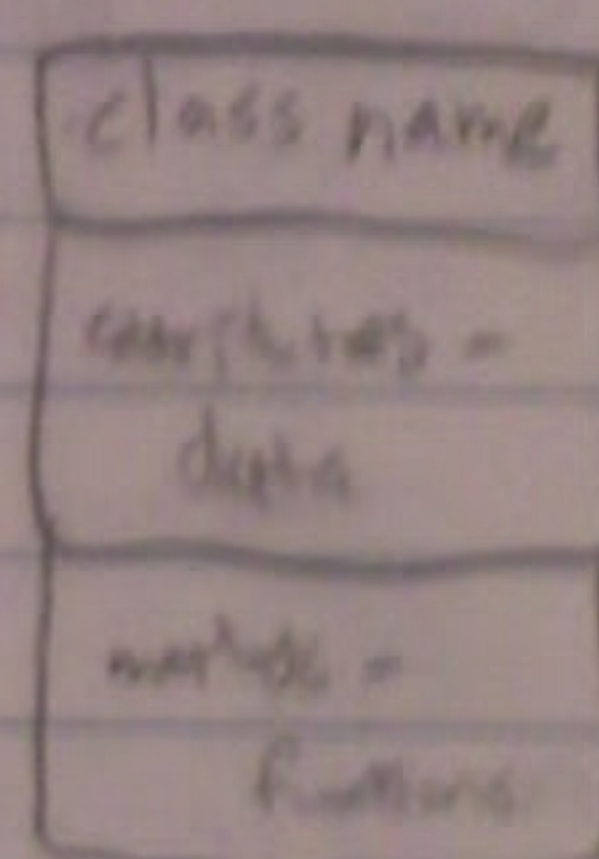
class = ex.

No hand drawn UML

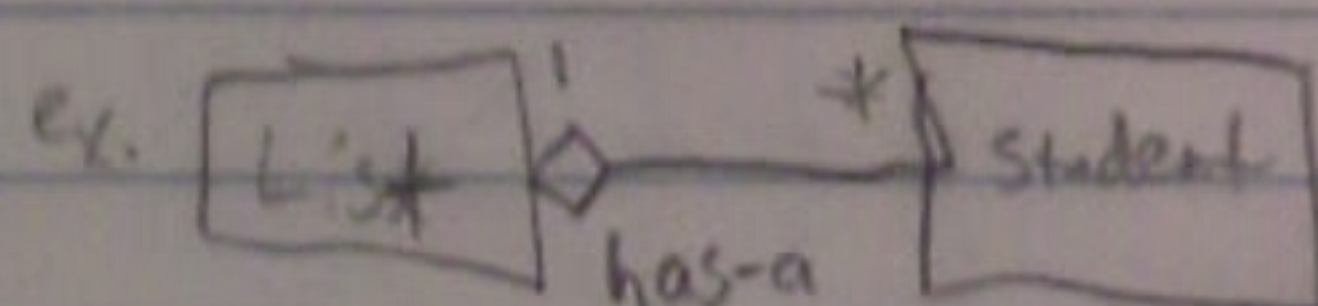
USE VIZIO for UML design

Zander only cares for composition,
inheritance.

real UML (might not be in 343)

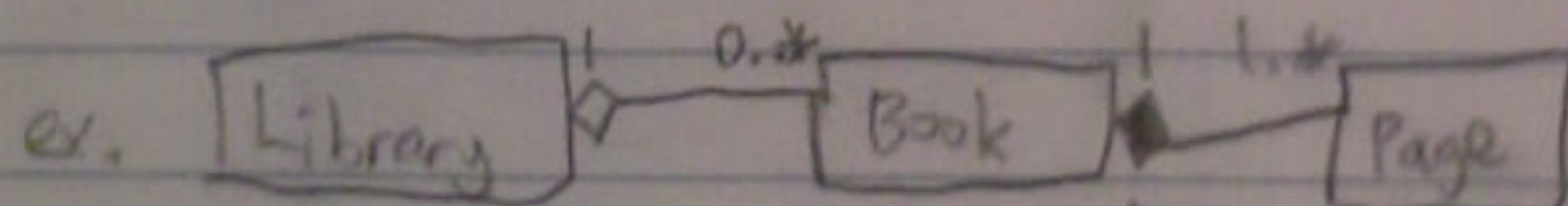


← not used



Composition (aggregation in Vizio)

1 List object will have many Student objects

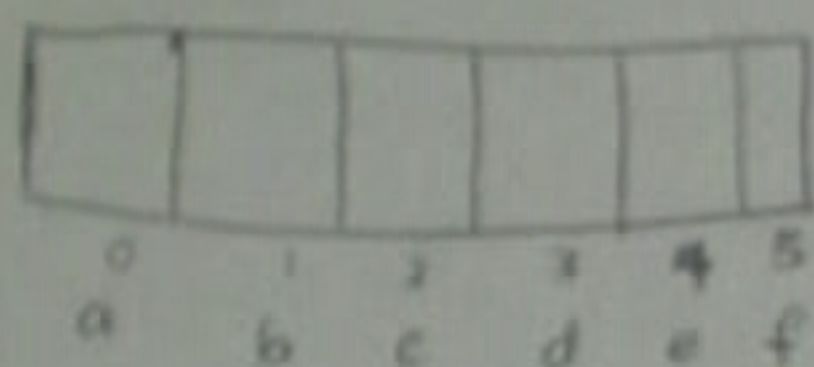


can exist if it
doesn't have any

can not exist if it doesn't
have any.

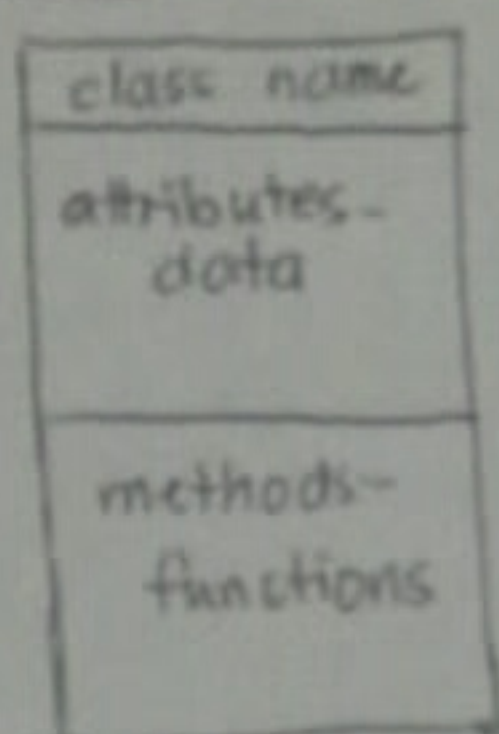
→ Cont.


```
char letter, [e]
int index = letter - 'a';
```



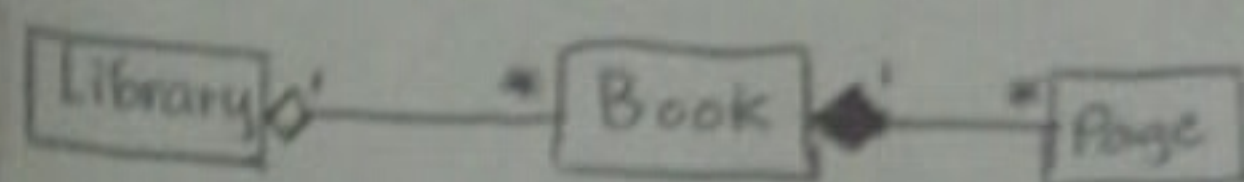
* This is a basic Hash Table

2/8/2014
CSS 343

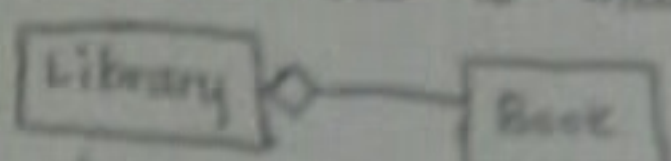


UML

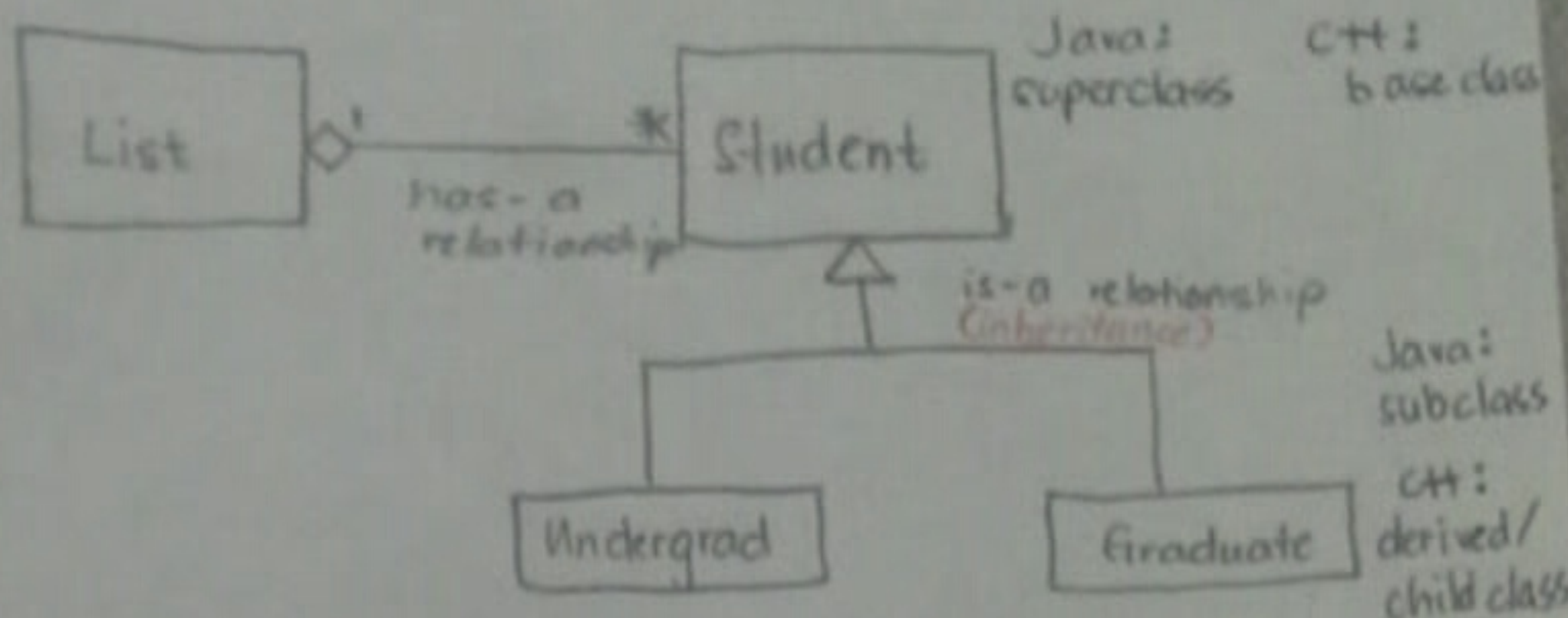
rectangle → class
diamond → composition / aggregation
L open ○ can exist if nothing
L closed ◆ can't exist if nothing
* many
1 one
hollow arrow (→) → inheritance



- if no pages, not a book
- if no book, can still be library
- 1 * : one to many



↳ Library has a book

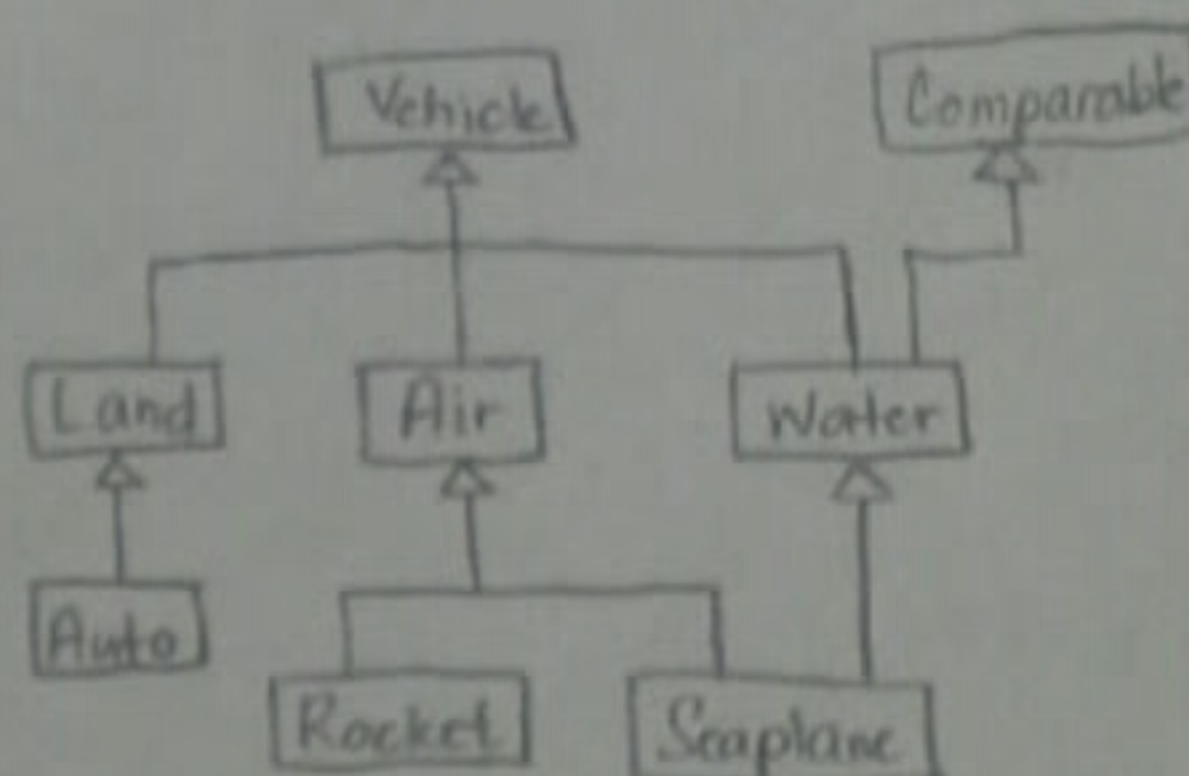


Nouns

table → manager
card
deck
player
dealer
hand
bet
pot

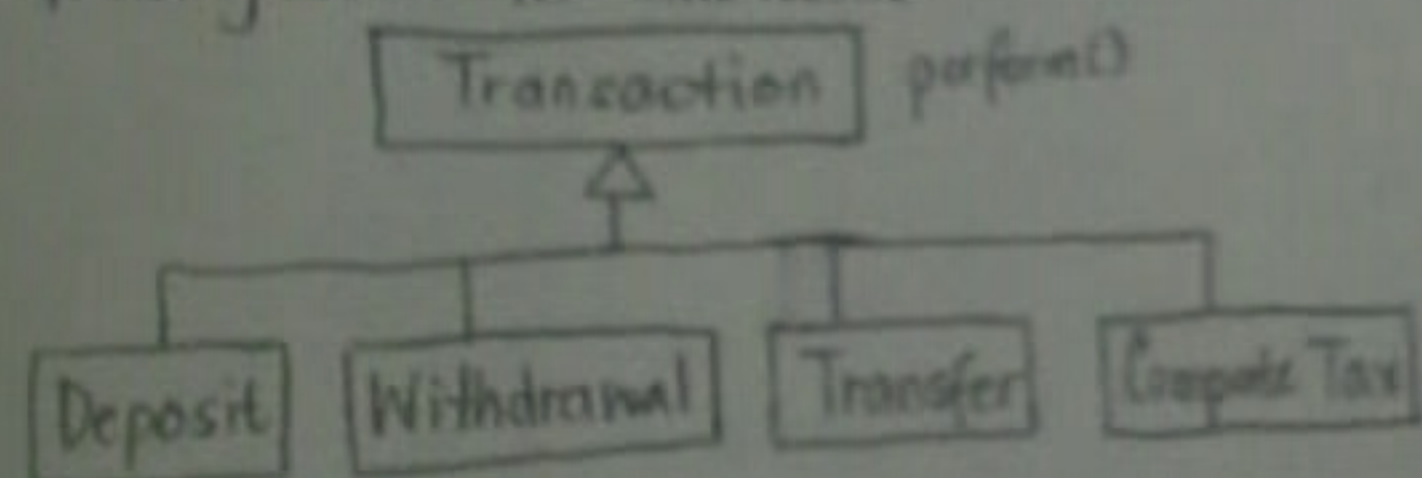
Verbs

shuffle
bet/raise/call
deal
fold



Lab 5 Design

* using verbs for inheritance:



Open-Closed principle → open for extension, closed for modification

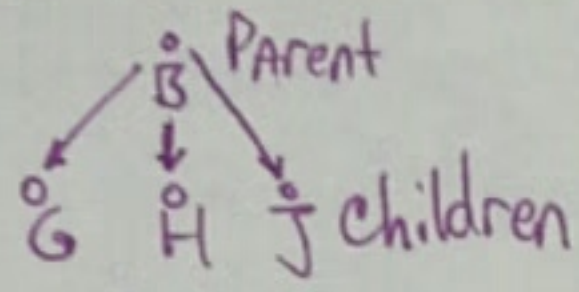
Polymorphism → doing the same tasks
different ways

@ - is not object-oriented
- structured language

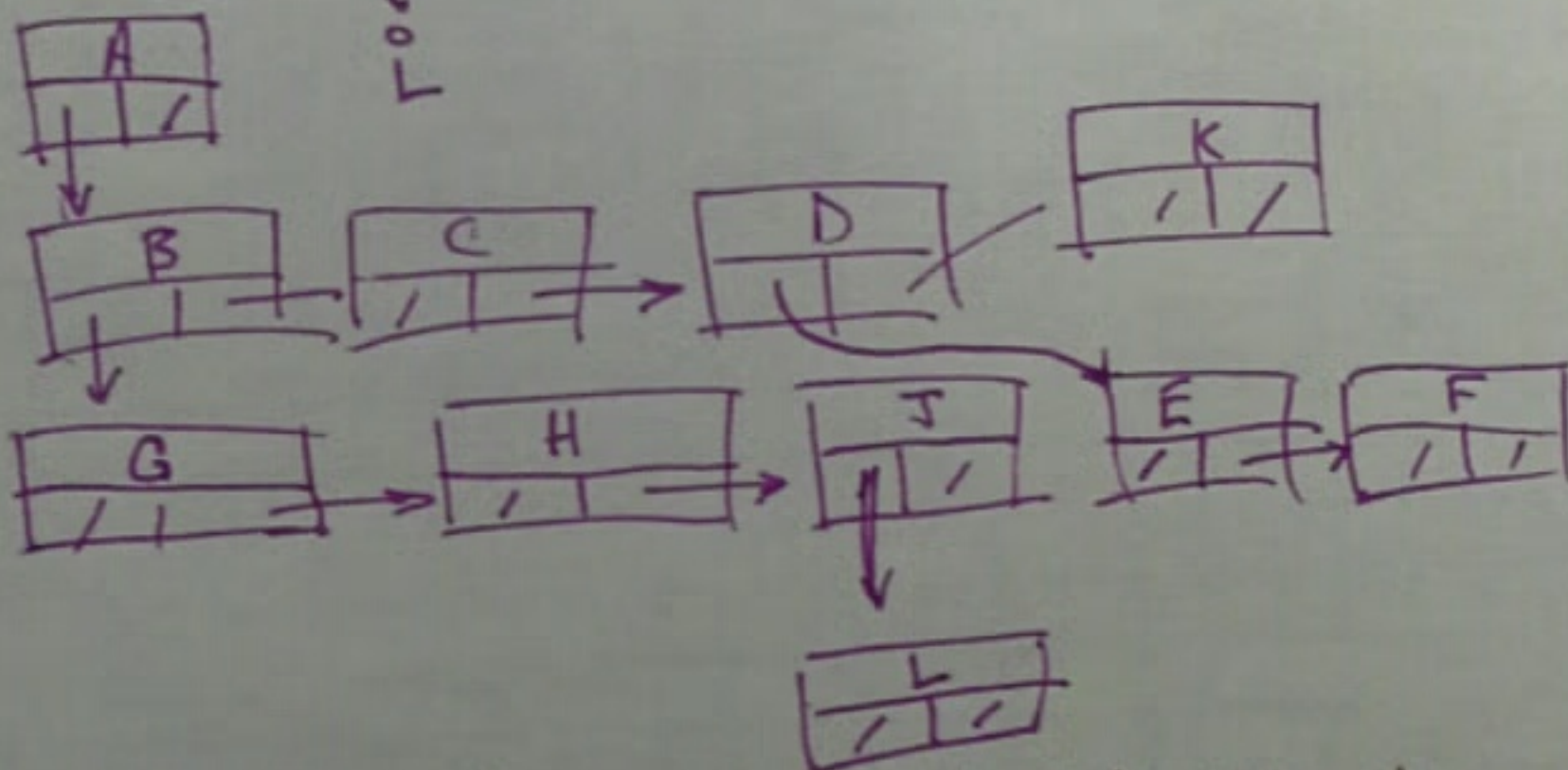
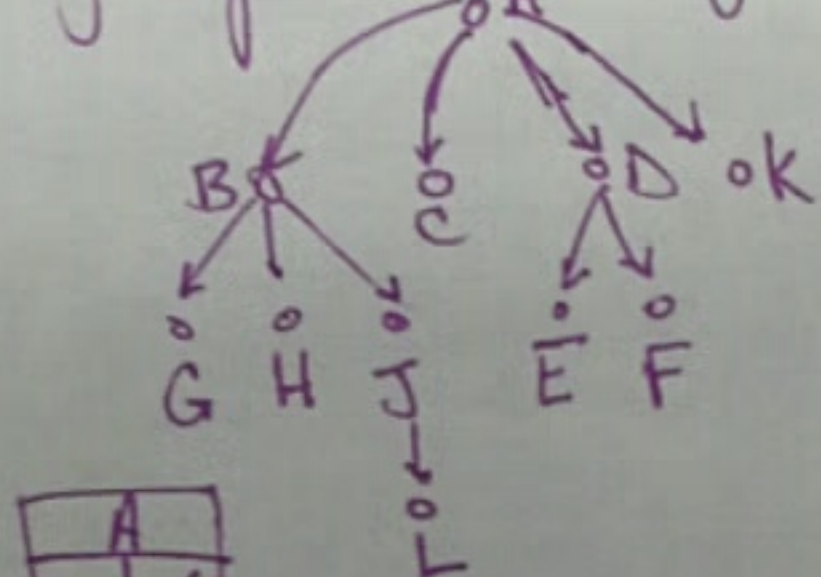
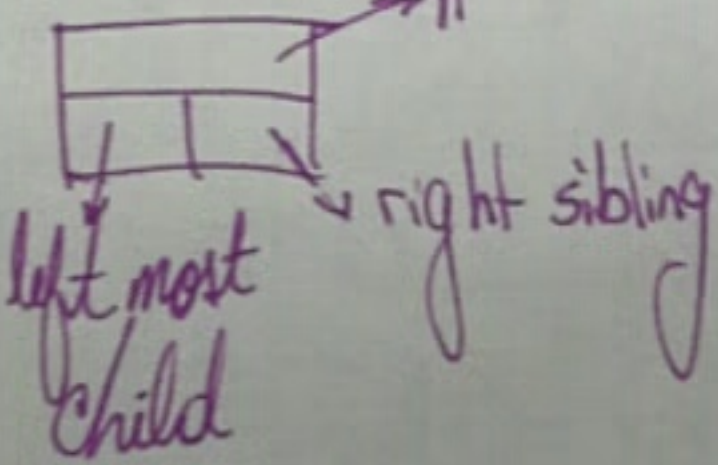
○ Bubblesort for comment how-to. What code is doing → not how.
 end column comments → keep it short

1/8/2014

Trees: A single node is a tree.
 G H J



Common implementation of a general tree - leftmost child, right sibling



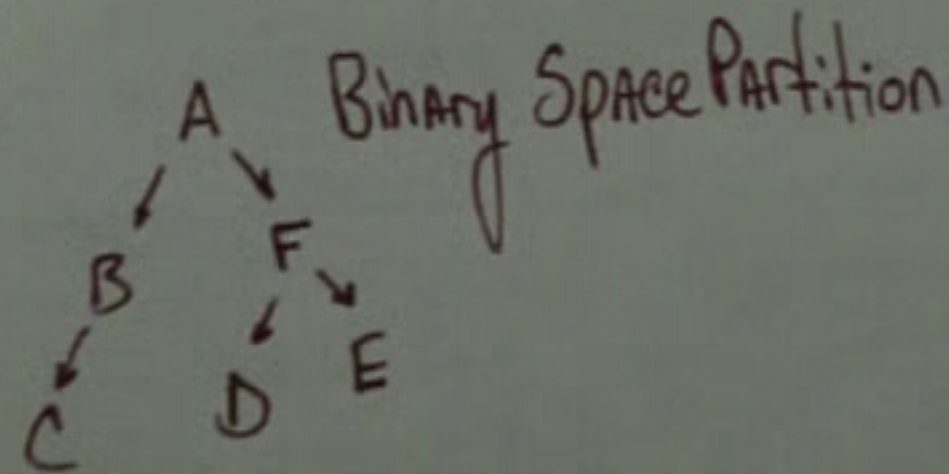
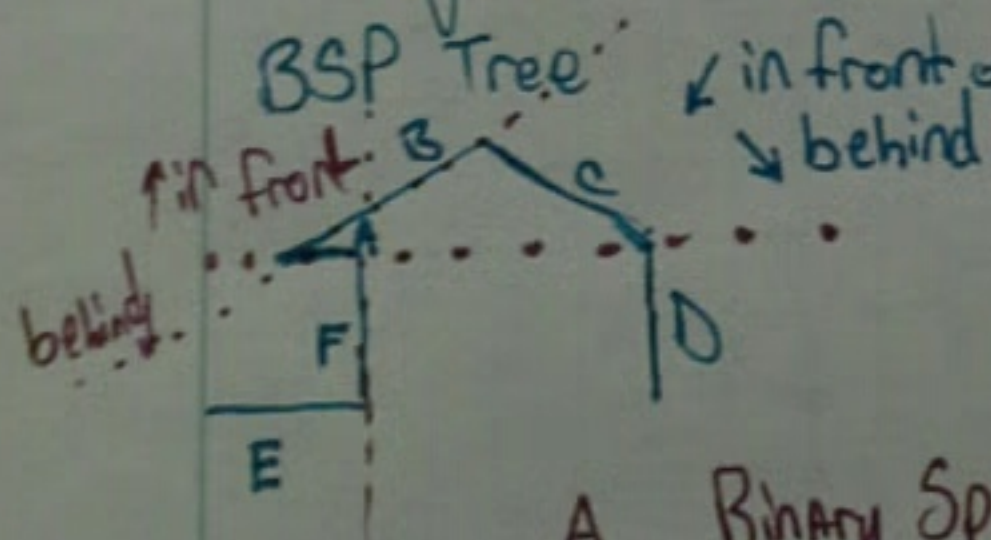
□ Expression Tree

□ BSP Tree

□ Huffman encoding algorithm

□ Binary Search Tree

Encoding - Find a uniquely encoded and decoded code s.t. chars w/ higher freq have shorter encodings



ADT Poly
 data structure → Array
 Binary Search Tree
 - finite ordered collection

Hashing notes

$N = \#$ of items in table

$B = \#$ of buckets, $\#$ of elements in array

Open Hashing / separate chaining



$$O(1 + \frac{N}{B}) = O(\frac{N}{B})$$



Closed Hashing / open addressing

First probe: $\text{Prob}(\text{1st collision}) = \frac{N}{B}$

Second Probe: $\text{Prob}(\text{2nd Collision}) = \frac{N}{B} \times \frac{N-1}{B-1}$

$$\text{Prob}(i^{\text{th}} \text{ Collision}) = \frac{N}{B} \times \frac{N-1}{B-1} \times \dots \times \frac{N-i+1}{B-i+1}$$

$$\approx O(\frac{N^i}{B^i}) = O(\frac{N}{B})^i$$

Average $\#$ of probes: $1 + \sum_{i=1}^{\infty} \left(\frac{N}{B}\right)^i = \frac{N}{B-N} = \frac{1}{1-(\frac{N}{B})}$

when found $\nearrow$

sum over all i
of at least i collisions

α called load factor

early binding - compile time

late binding $\rightarrow$ run time

start with fiction
get out put

move on once finished
add next content

add Patron

Add Actions

Polymorphism d.vptr



vtable holds addresses

of any virtual functions

every class with virtual has
a vtable.

Polymorphism

class B

public:

```
virtual char f() { return 'B'; }
char g() { return 'B'; }
char testF() { return f(); }
char testG() { return g(); }
```

class D: public B

public:

```
virtual char f() { return 'D'; }
char g() { return 'D'; }
```

Main

```
D d;
cout << d.testF() << " " << B
d.testG() << endl;
cout << d.g();  $\rightarrow$  D
```


- Changed content factory
- action factory
- gonna override display & actions ∈ History.
↳ History has-A actions

• links to composition & graphic

• operator overloads all pure virtual

↳ Super class

NodeData → KindA Like JAVA's object

No NodeData → use content as NodeData

↳ setData - override & content
↳ give ioStream

History holds actions

pass actions as ptr.

Language - a set of strings

We're interested in expressiveness of languages

Compilers

Stream of chars

lexical analyzer

tokens

syntactical
semantic
analyzer
"Parser"

Chomsky Hierarchy

Tokens

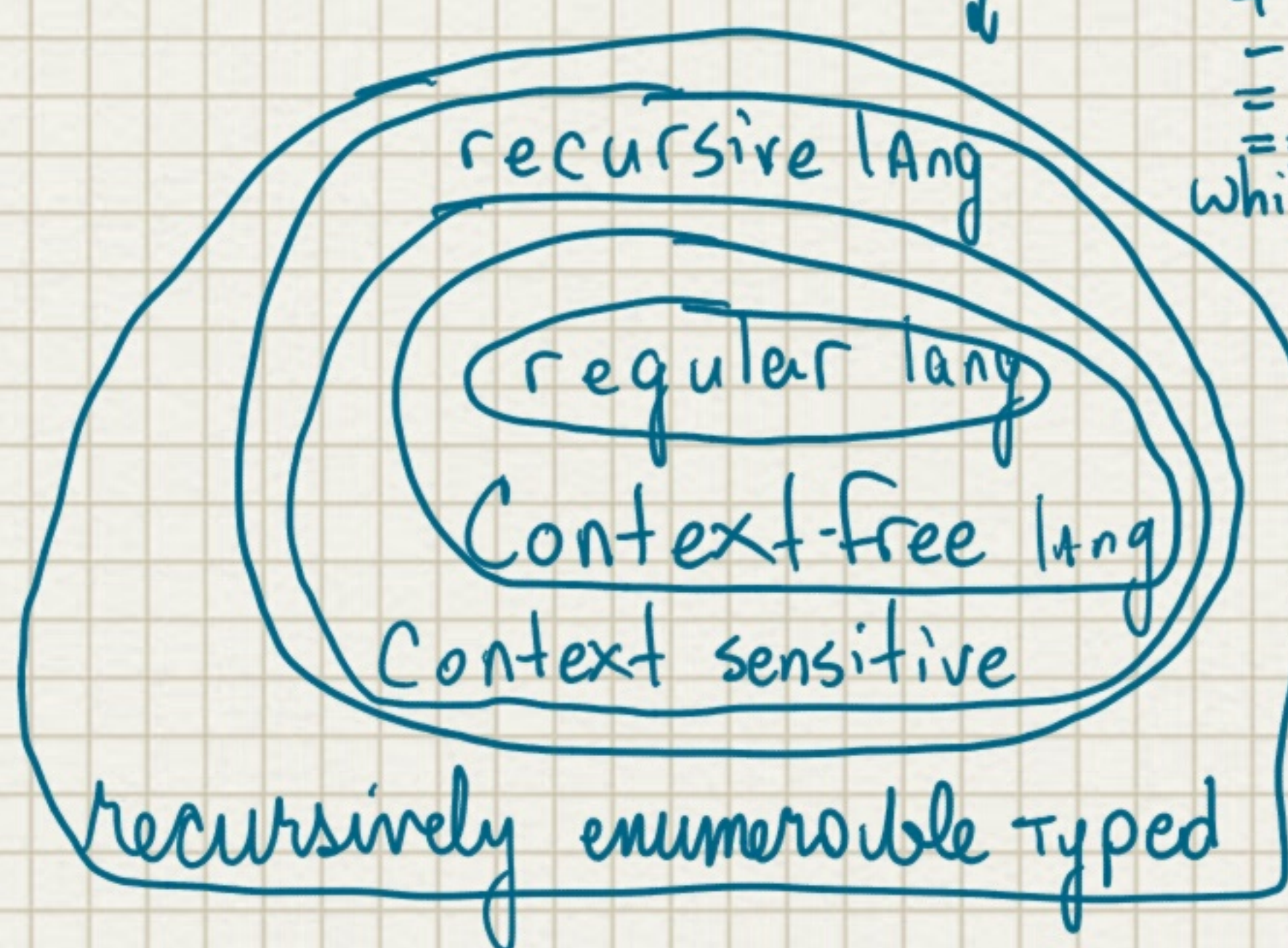
(
+
=
while

identifier: book

number: 50,100.27

string literal "m"

char



even a's and b's

$(a|b|(ab|ba)(a|b)^*(ab|ba))^*$

even a's
 $(ab^*ab)^*$

odd a's $(ab^*ab)^*ab^*$

$\Sigma = \{0, 1, +, -\}$

$L_1 = \{S \mid S \text{ is an unsigned binary number}\}$

$(+|-|\lambda)(0|1)^*(0|1)^*$

Regular expressions

Have alphabet, often called Σ

1. empty string is reg expr (λ, ϵ)

2. Every char in alphabet is a reg expr

3. If x, y are reg exprs, then so are

highest precedence

(x)

x^* Kleene star, repetition
(zero or more)

xy concatenation

$x|y$ OR, sometimes see $x+y$

lowest

Housekeeping

1. Do I generate them all?
2. Do I generate any string I am not supposed to generate?

$$L_1 = \{s \mid s = na + mb, n, m \geq 0\} \quad L_2 = \{s \mid na + nb, n \geq 0\}$$

"a*b"

$\Sigma = \{a, b, 1, 2, -\}$ C++ like identifier

1. must start with letter, or underscore
2. Then any letters, digits, underscores

$(a|b|_)(a|b|1|2|_)^*$

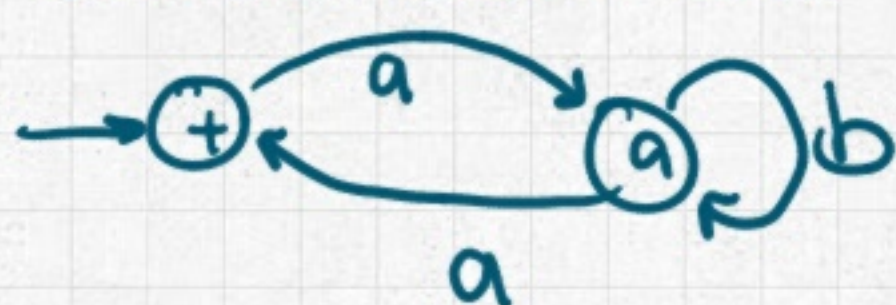
Ada-like identifier

1. must start with letter
2. Then letters, digits, underscores

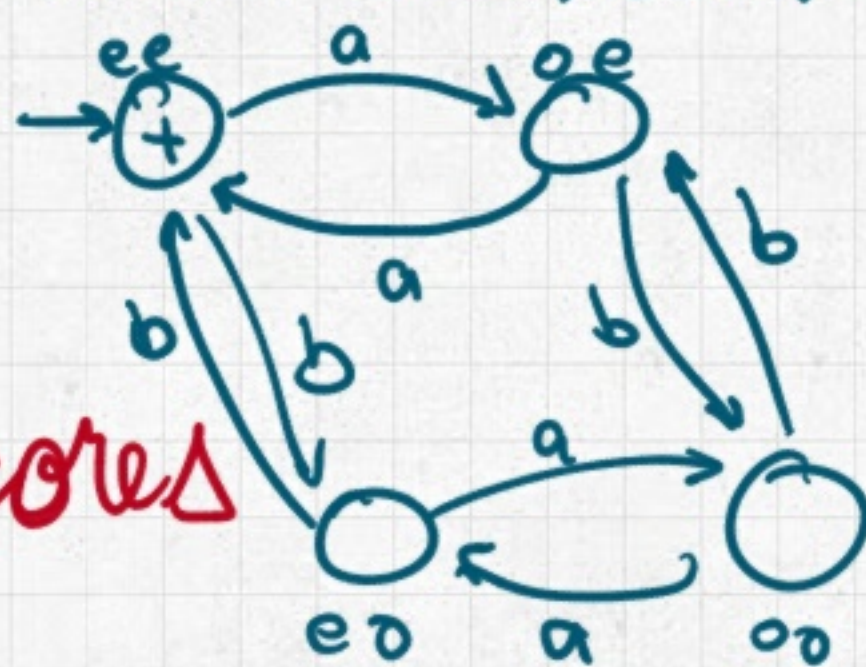
except cannot have two consecutive underscores

nope $(a|b)((a|b|1|2|_)(a|b|1|2)(a|b|1|2|_))^*$

$\Sigma = \{a, b\}$ Even a's



Even a's, Even b's

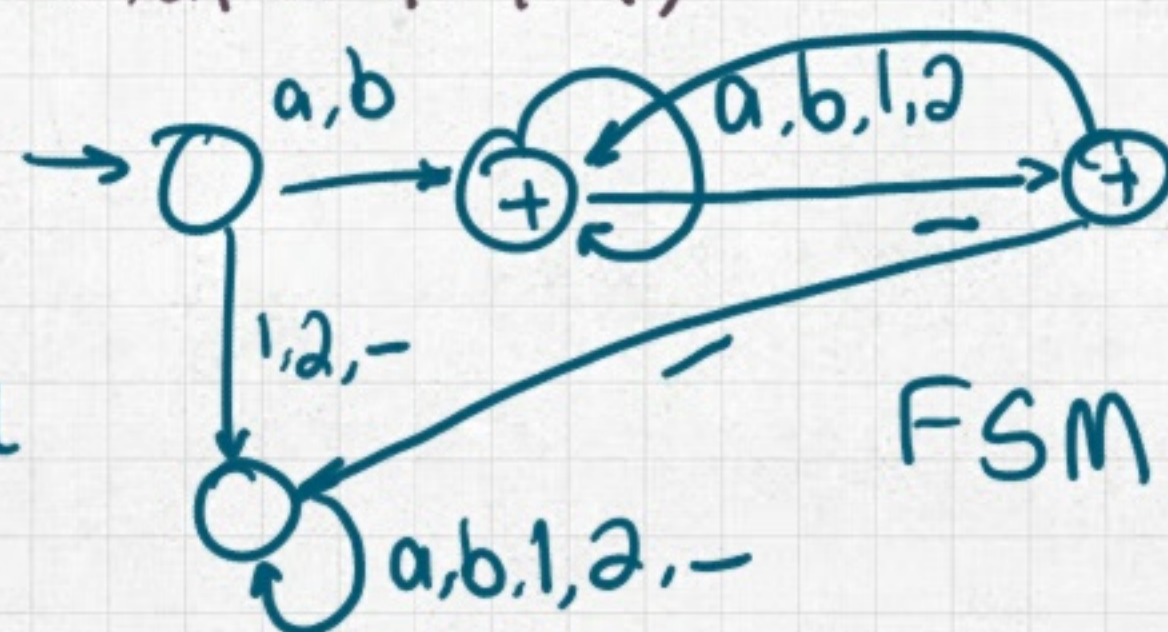


no $(a|b)((a|b|1|2|_)(a|b|1|2))^*$

nope $(a|b)((a|b|1|2|_)(a|b|1|2)^*)(a|b|1|2|_)$

Still no $(a|b)((1|a|b|1|2|_)(1|a|b|1|2) | (1|a|b|1|2)(1|a|b|1|2|_))^*$

yes $(a|b)((a|b|1|2|_)(a|b|1|2))^*(-|1)$



Finite State Machine $\leftrightarrow$ Finite Automata

Deterministic Finite State Machine

Kleene's Theorem

$\forall$ Regular Language

1. recognized by RegEx
2. FSM recognizes
 - i. set of states
 - ii. start state $\ominus \rightarrow \circ$
 - iii. final state $\oplus \rightarrow \odot$
 - iv. alphabet - input
 - v. transition function

Milan's thoughts

- 1/ finish display of each type of content
- Write display Library
 - Columns of data → aligned!
- History has many actions
 - ↳ display() $\forall$ actions $\in$ History → overwrite
 - ↳ implement check out()
 - " return()
 - " display Patron history

Stretch Goals

Format names
Format headers
A.V.L. our trees

To Do: Abstract out
content factory from content.

North

As Waltz

Danny

Horton

I hate

Martha

Trehorn

Gentle

If you

Smiley

Old

The Sea

Roberts

Penelope

lexical analysis

tokens

parser

YACC
Bison

Context-Free Grammar Write BNF

1. set of terminals (characters, tokens)
2. set of non-terminals
3. starting non-terminal
4. set of production rules → Looks like

one nonterminal → terminal and nonterminal

$\langle \text{sentence} \rangle \rightarrow \langle \text{noun-phrase} \rangle \langle \text{verb} \rangle \langle \text{noun-phrase} \rangle$

$P \langle \text{noun-phrase} \rangle \rightarrow \langle \text{noun} \rangle$ equivalent $\leftrightarrow [\langle \text{noun-phrase} \rangle \rightarrow \langle \text{noun} \rangle | \langle \text{noun-phrase} \rangle \rightarrow \langle \text{adj} \rangle \langle \text{n-p} \rangle]$

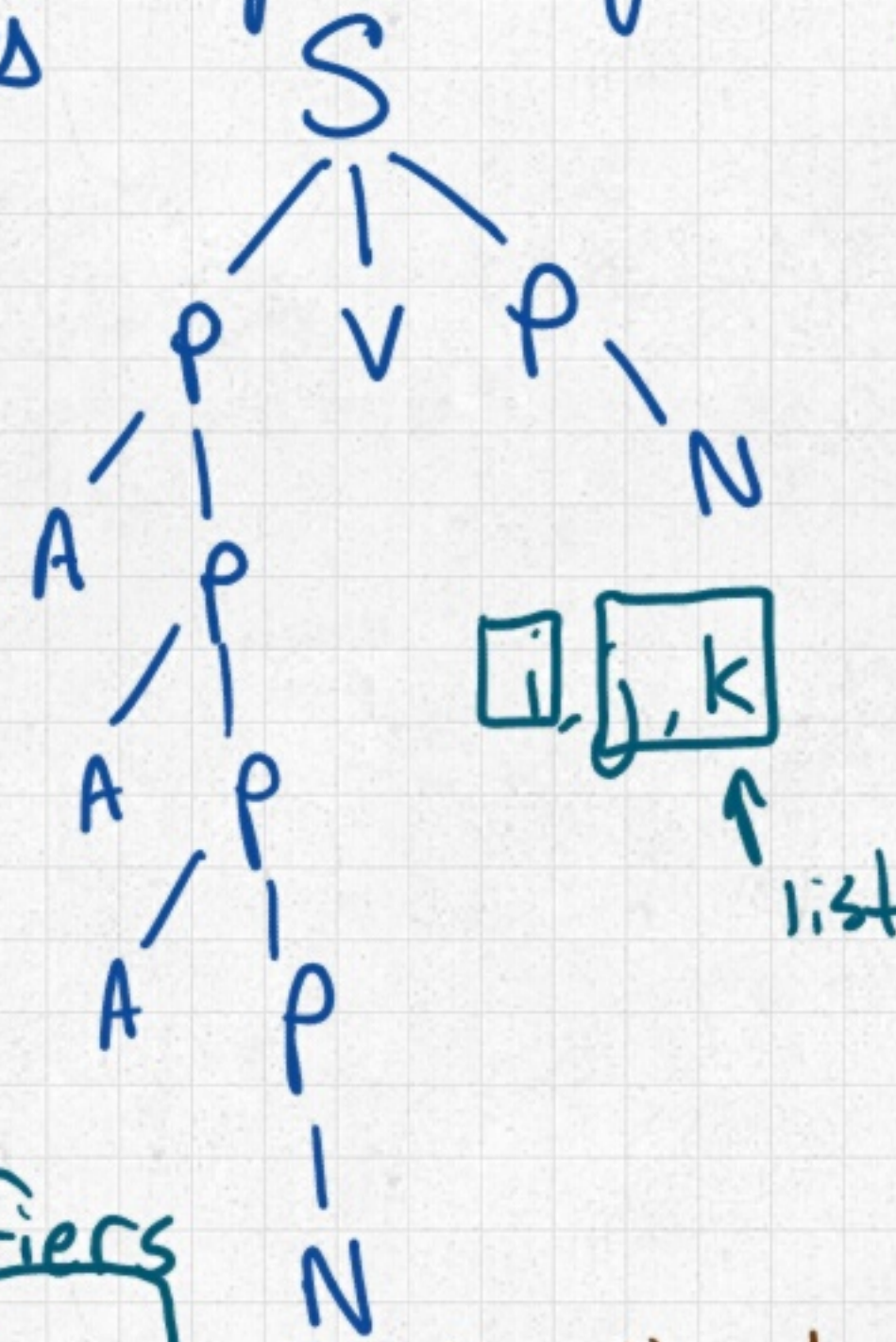
$P \langle \text{noun-phrase} \rangle \rightarrow \langle \text{adj} \rangle \langle \text{noun-phrase} \rangle$

$N \langle \text{noun} \rangle \rightarrow \text{cat} | \text{dog} | \text{fish} | \text{apple} | \text{ball}$

$V \langle \text{verb} \rangle \rightarrow \text{eat} | \text{pet} | \text{juggle}$

$A \langle \text{adj} \rangle \rightarrow \text{tasty} | \text{brilliant} | \text{red}$

Demonstrate generating string
2. Parse Trees



Derivation

1. $S \Rightarrow PVP$

$\Rightarrow APVP$

$\Rightarrow AAPVP$

$\Rightarrow AAAPVP$

$\Rightarrow AAANVP$

$\Rightarrow AAANVN$

`int i,j,k;
float x,y;
char a,b,c;
bool done;`

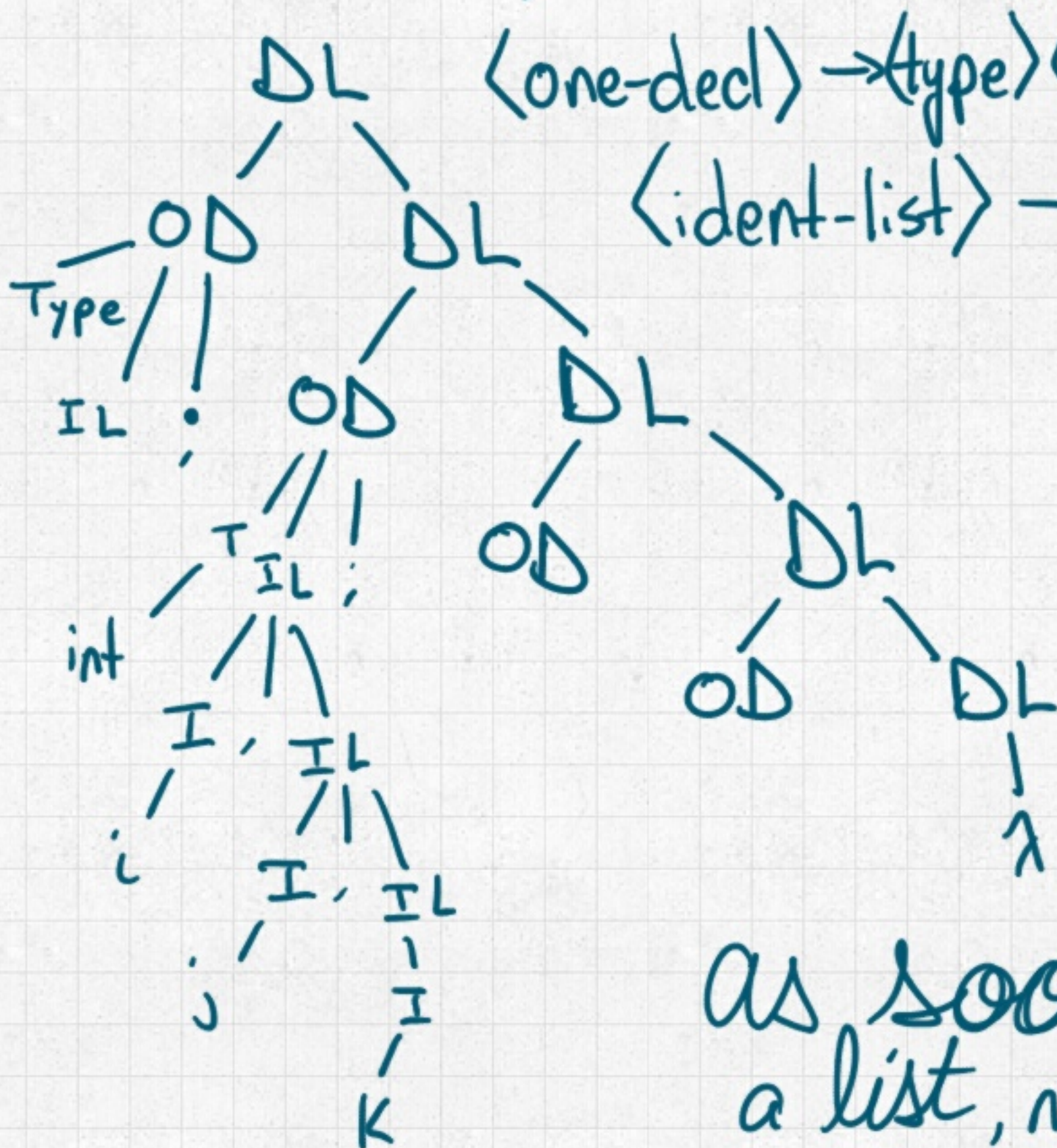
$\langle \text{decl-list} \rangle \rightarrow \lambda$

list $\xrightarrow{\text{make}}$ nonterminal

$\langle \text{decl-list} \rangle \rightarrow \lambda | \langle \text{one-decl} \rangle \langle \text{decl-list} \rangle$

$\langle \text{one-decl} \rangle \rightarrow \langle \text{type} \rangle \langle \text{ident-list} \rangle ;$

$\langle \text{ident-list} \rangle \rightarrow \langle \langle \text{identifier} \rangle \rangle$



Identifiers

while

>

<

=

==

{

}

identifier

number

literal string

Bad Arithmetic Expression grammar

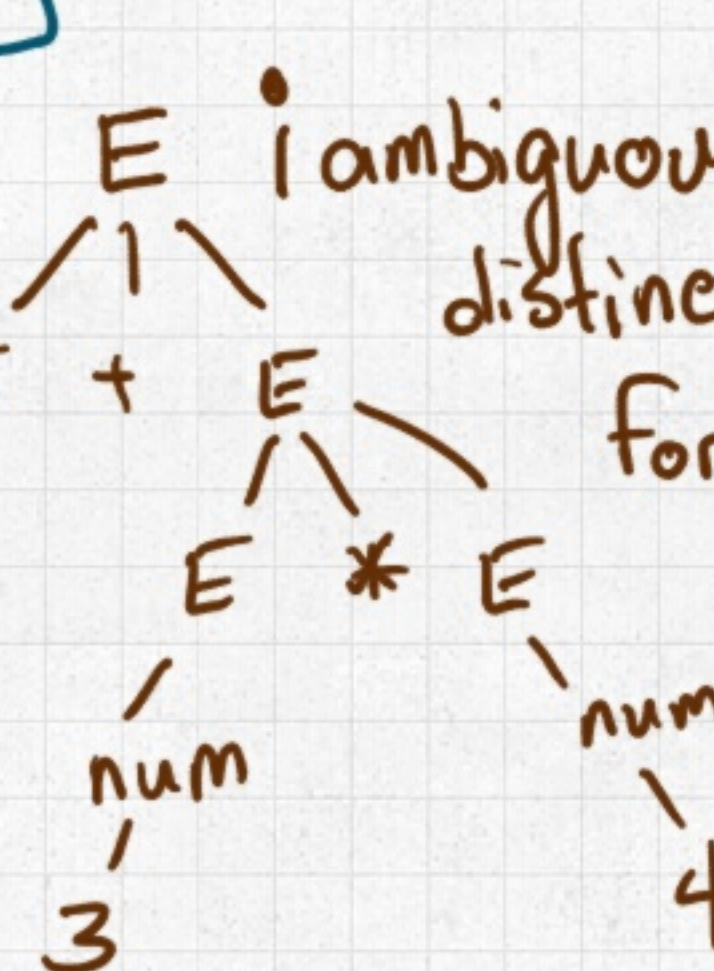
$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{expr} \rangle |$

$\langle \text{expr} \rangle - \langle \text{expr} \rangle$

$E \rightarrow E + E | E - E | E * E | E / E |$

$(E) | \langle \langle \text{number} \rangle \rangle | \langle \langle \text{ident} \rangle \rangle$

ambiguous $\leftrightarrow \exists > 1$ distinct parse tree for 1 same string.



Good Expression Grammar

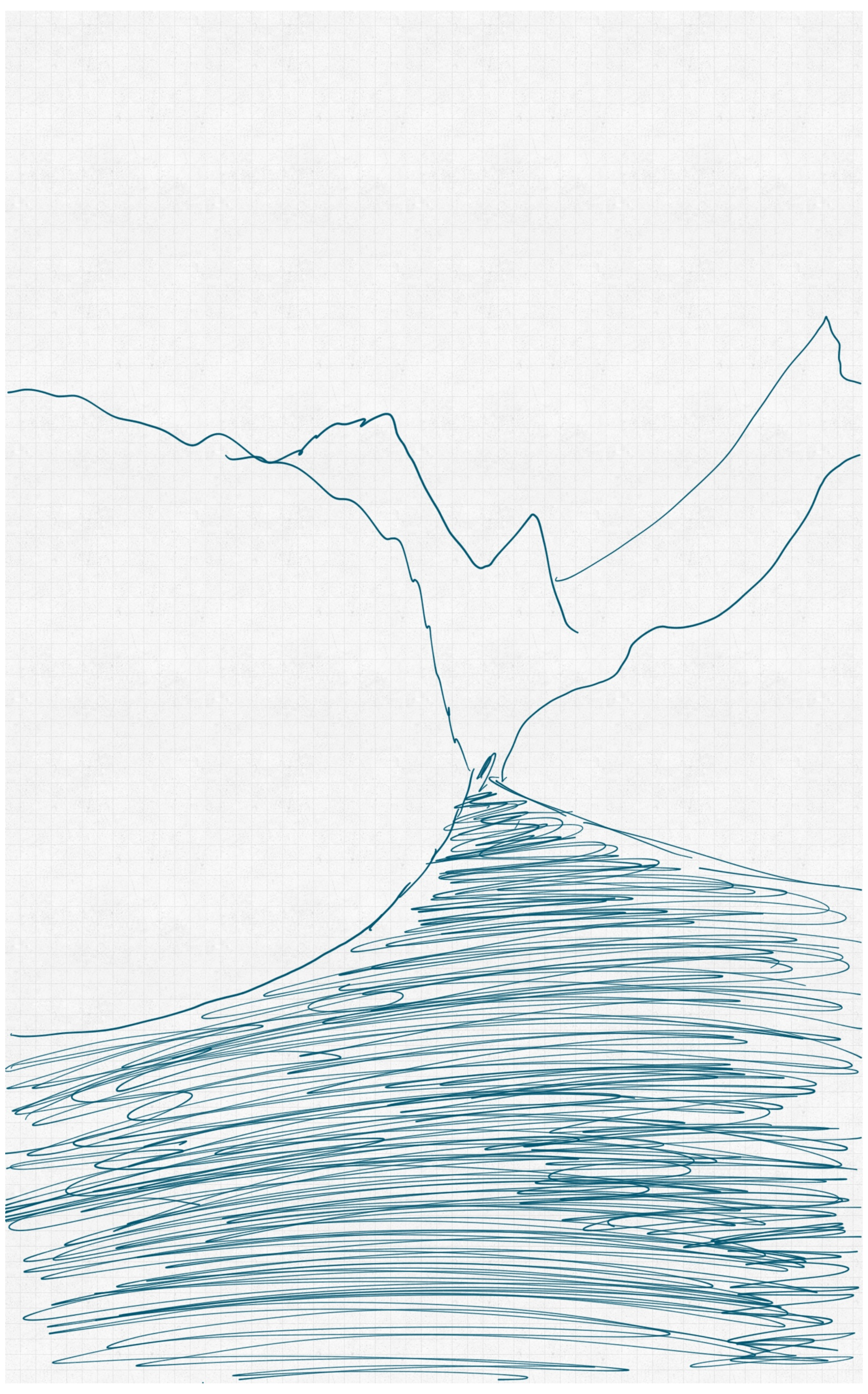
$E \rightarrow E + T | E - T | T$

$T \rightarrow T * F | T / F | F$

$F \rightarrow (E) | \langle \langle \text{number} \rangle \rangle | \langle \langle \text{ident} \rangle \rangle$

rule: haz identifier

→ comma, ident-list



Semi-comprehensive

↳ Focus on second half

↳ 8x11 Cheat sheet, Both sides

↳ ^{till} 10:30

Hashing → linear/Quadratic Probing

Complexities

inheritance/polymorphism → design

Open/Close principle

↳ code: read/write (Polymorphism)

at least one (2?) focused on Poly/inher^{it}

implementation under the hood

Regular Expression

↓ Languages ↑ Turing Machines (about it)^{stuff}

↳ "Put it all together"

↳ Factory Pattern/Command Parent

Transaction

Checkout return